

These are some review questions to test your understanding of the material. Some of these questions may appear on an exam.

1 Priority Queues and Heaps

Note: pseudocode for merging leftist heaps is given on page 3.

- 1.1 Define “priority queue.”
- 1.2 Define “binary heap.”
- 1.3 Define “null path length” in a binary tree.
- 1.4 Define “leftist binary tree.”
- 1.5 Define “leftist heap.”
- 1.6 Insertion and deletion in a binary heap is in $O(\lg n)$ on average. Explain why this is so.
- 1.7 Finding the minimum element in a binary heap (a min-heap) is in $O(1)$ worst case. Explain why this is so.
- 1.8 Prove that the largest element in a min binary heap is a leaf.
- 1.9 For a min binary heap of N elements, what is the range of indices in which the maximum element will be found?
- 1.10 Describe, in English, an algorithm to find the largest element in a binary min-heap. What is the asymptotic worst case performance of your algorithm?
- 1.11 The array representing a binary heap contains the following elements in the order 2, 8, 3, 10, 16, 7, 18, 13, 15. Show the order that results at **each** stage of inserting the element 4.
- 1.12 The array representing a binary heap contains the following elements in the order 2, 8, 3, 10, 16, 7, 18, 13, 15. Show the order that results at **each** stage of deleting the minimum element.
- 1.13 Describe, in English, the process for constructing a binary heap from a given set of initial values (*i.e.*, heapify an array of elements).
- 1.14 Construct a binary heap using the initial values 18, 2, 13, 10, 15, 3, 7, 16, 8. Show the heap at **each** stage of construction.
- 1.15 Given a drawing of a binary tree, state if it is a leftist tree and if it is a leftist heap. Give reasons.
- 1.16 Prove that any complete binary tree is a leftist tree.
- 1.17 Prove: for any leftist tree having N vertices, the number of vertices, R , on the rightmost path to a non-full vertex is given by

$$R \leq \lceil \lg(N + 1) \rceil$$

- 1.18 Describe how to do `findMin` in a leftist heap.
- 1.19 Using the `merge` (*aka* `meld`) operation, describe how to do the `insert` and `deleteMin` operations for leftist heaps.
- 1.20 Describe a method to construct a leftist heap from a given set of values. Your code must construct the leftist heap of N elements in $O(N)$ time (**not** in $O(N \lg N)$).
- 1.21 Given drawings of two leftist heaps `H1` and `H2`, draw the leftist heap that results from the operation `merge(H1,H2)`.

Merging Leftist Heaps

Here is pseudo-code for merging two leftist heaps. This is a destructive operation.

```
template <class Comparable>
Comparable & rootValue(LHeap<Comparable> & H)
{
    return the value stored at the root of H
}

template <class Comparable>
LHeap<Comparable> & leftChild(LHeap<Comparable> & H)
{
    return the left child of the root of H
}

template <class Comparable>
int leftLength(LHeap<Comparable> & H)
{
    return the null path length of the left child of the root
        of H
}

template <class Comparable>
LHeap<Comparable> & Merge(LHeap<Comparable> & H1,
                        LHeap<Comparable> & H2)
{
    if ( (H1 is empty) ||
        (rootValue(H1) > rootValue(H2)) )
        // make sure min value in H1 is not larger than
        // min value in H2
        swap(H1, H2);

    if (H1 is empty) // if is, then both H1 and H2 are empty
        return H1; // return an empty LHeap (or NULL)

    // recursive call to Merge
    replace rightChild(H1) by Merge(rightChild(H1), H2);

    // make sure H1 is still leftist
    if (leftLength(H1) < rightLength(H1))
        swap(left(H1), right(H1));

    return H1; // return the modified H1
}
```