

Homework #1

- Should be turned in by now!
- [Roughly] 25 people say marshmallow, 7 people said Jelly, 3 said either, 2 said both
- Some answers to questions and concerns:
 - The class is about math **and** computer science
 - We will learn math, and apply it to CS
 - You won't need to know a lot of math ahead of time
 - You won't need to know programming either
 - Getting an A and understanding concepts is up to you
 - If you don't understand something, ask questions on Piazza, email, and office hours
 - I will let you know what exams and quizzes will be like
 - There may or may not be donuts

Some Additional Notes

- Homework #2 will be released today
 - Will be due next Thursday
- Piazza now contains slides and homeworks
- **Extra credit:** Come to my office hours during the next week; you don't need a question, just show up
- *Lemma:* “proven statement used as a stepping-stone toward the proof of another statement” – basically a small theorem that is used in another proof

Lightning Review



Lightning Review (Propositions)



Which of these are propositions?

- Do not pass go.
- What time is it?
- $2 + 3 = 5$
- $4 + x = 5$
- The moon is made of green cheese.
- Answer this question.
- You answered that question.

Lightning Review (Logic $\rightarrow$ English)



Convert to English, assuming p and q are propositions

p : I bought a lottery ticket this week

q : I won the million dollar jackpot

(a) $p \rightarrow q$

(b) $p \wedge q$

(c) $\neg p \vee (p \wedge q)$

(Hint: Order of operations)

(d) $p \leftrightarrow q$

Lightning Review (English $\rightarrow$ Logic)



Write using logical connectives, assuming:

p : It is below freezing

q : It is snowing

- (a)** If it is snowing, then it is freezing
- (b)** That it is below freezing is necessary and sufficient for it to be snowing
- (c)** It can be snowing or below freezing (or both)
- (d)** It is not below freezing and it is not snowing

Lightning Review (Conditionals)



Determine if these propositions are true

- If $1 + 1 = 2$, then $2 + 2 = 5$
- If $1 + 1 = 3$, then $2 + 2 = 5$
- $1 + 1 = 2$ if and only if $2 + 3 = 4$
- $1 + 1 = 3$ if and only if monkeys can fly
- If there is a pop quiz, then we will pass

Lightning Review (Truth Tables)



Construct a truth table for:

- $(p \oplus q) \vee (p \oplus \neg q)$
- $(p \rightarrow q) \wedge (\neg p \rightarrow r)$

Label the following as either a *tautology*, *contradiction*, or *contingency*?

$p \wedge \neg p$	$(p \vee q) \rightarrow (p \vee q)$
$p \wedge q$	$(p \vee q) \wedge (\neg p \wedge \neg q)$
$p \vee \neg p$	$p \wedge p$

CMSC 203: Lecture 3

Some Reminders

Reminder of why this matters

- What is the use of logic and proofs in the field of:
 - computer engineering?
 - computer programming?
 - artificial intelligence?
 - mathematics?
 - game development?

Fun Equivalences

- $p \rightarrow q \equiv \neg p \vee q$
- $p \leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$

- **Identity Laws**

- $p \wedge T \equiv p$
- $p \vee F \equiv p$

- **Domination Laws**

- $p \vee T \equiv T$
- $p \wedge F \equiv F$

- **Negation Laws**

- $p \vee \neg p \equiv T$
- $p \wedge \neg p \equiv F$

- **Idempotent Laws**

- $p \vee p \equiv p$
- $p \wedge p \equiv p$

- **Commutative Laws**

- $p \vee q \equiv q \vee p$
- $p \wedge q \equiv q \wedge p$

More Fun Equivalences

- **Absorption Law**

- $p \vee (p \wedge q) \equiv p$
- $p \wedge (p \vee q) \equiv p$

- **Distributive Law**

- $p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$
- $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$

- **De Morgan's Law**

- $\neg(p \wedge q) \equiv \neg p \vee \neg q$
- $\neg(p \vee q) \equiv \neg p \wedge \neg q$

Why Equivalences?

- How can equivalences be used in the following:
 - in writing a proof?
 - in programming?
 - in writing an artificial intelligence?

Proving Equivalence

Assume our robot knows $\neg p \vee \neg q$ is true

Can it know whether $\neg(p \wedge (\neg p \vee q))$ is true?

- It can not figure out the values of p or q
- It must prove that $\neg p \vee \neg q \equiv \neg(p \wedge (\neg p \vee q))$
- If it is equivalent, we don't need the values of p or q
- Equivalence can be proven with a proof table (slow!)
- Or we may attempt to use our laws to reform
 - (Think back to *solve for x* in algebra!)

CMSC 203: Lecture 3

Predicate Logic

Predicate Logic

- Problems with Propositional Logic
 - We can not express meaning of all statements in mathematics and natural language
- Example:
 - Assume every computer connected to the school network is functioning properly
 - Determine if MATH3 is functioning properly
- Predicate logic is more powerful than propositional logic

What is in Predicate Logic?

“Computer x is functioning properly”

- **Variables**

- x is the variable (a la computer programs / algebra)
- Subject of the statement

- **Predicate**

- *“is functioning properly”* is the predicate
- Properties the subject can have
- Think of them as **functions**
 - You may even express them as functions

Predicates

- Predicates can be expressed as $P(x)$
 - Similar to functions, x is some input
 - Once x is assigned a value, $P(x)$ becomes a proposition (*recall what a proposition is*)
- Example:
 - $P(x)$ denotes " $x > 3$ "
 - What is the value of $P(4)$ and $P(2)$?
 - $R(x, y, z)$ denotes " x hates y , but y loves z "
 - What is the truth value of $R(\text{Tybalt}, \text{Romeo}, \text{Juliet})$

Quantification

- **Quantifiers**

- Designate a *range* in which a predicate is true
- Two kinds we will look at:
 - **Universal** – Predicate is true under every element
 - Eg: “*Everybody* loves Raymond”
 - **Existential** – Predicate is true for *at least* one
 - Eg: “*Somebody* is President”

Universal Quantifier

- Asserts predicate is true for all variables in **domain**
- $\forall x P(x)$ denotes the universal quantification for $P(x)$

Read: “for all/every x $P(x)$ ”

- **True** when $P(x)$ is true for *every* x
- **False** when there is an x for which $P(x)$ is false

Existential Quantifier

- Asserts an element exists in **domain** for which predicate is true
- $\exists x P(x)$ denotes the existential quantification for $P(x)$

Read: "There exists some x such that $P(x)$ "

- **True** when there is an x in which $P(x)$ is true
- **False** when $P(x)$ is false for every x

Quantifier Practice

- $\exists x \text{ 203Student}(x) \wedge \text{202Student}(x)$
- $\forall x \text{ 203Student}(x) \rightarrow \text{Sophomore}(x)$
- $\forall x \text{ Senior}(x) \wedge \text{Tall}(x)$
- Every gardener likes the sun.
- All purple mushrooms are poisonous.
- No purple mushroom is poisonous.

De Morgan's Laws (again!)

- $\neg \exists x P(x) \equiv \forall x \neg P(x)$
- $\neg \forall x P(x) \equiv \exists x \neg P(x)$
- There does not exist an element that... $\equiv$ For all elements, it is not true that...



Nested Quantifiers

- Quantifiers can have *scope*
- We can stack / nest our quantifiers for more flexibility
- For example: $\forall x \exists y \equiv$ “For all x , there exists a y ...”
- **Order is important!**
- Think of them as nested “loops”

Nested Quantifier Examples

- $\exists x \exists y \text{ Montague}(x) \wedge \text{Capulet}(y) \rightarrow \text{Enemies}(x, y)$
- $\forall x \forall y (x + y = y + x)$
- $\forall x \exists y \text{ Father}(x, y) \wedge \text{Male}(y)$
- $\exists x \forall y (\text{Student}(x) \wedge \text{Hard}(y)) \rightarrow \text{Failed}(x, y)$

Inference

- Special case of a mathematical proof
- Show some sentence p is true, given a set of **premises**
 - p is called the **conclusion**
- Our **argument** is **valid** if the truth of the conclusion follows the premises
- For example, given *premises* p and $p \rightarrow q$, we can reach the *valid conclusion* q

Rules of Inference

$$\begin{array}{l} p \\ p \rightarrow q \\ \hline \therefore q \end{array}$$

Modus
ponens

$$\begin{array}{l} p \rightarrow q \\ q \rightarrow r \\ \hline \therefore p \rightarrow r \end{array}$$

Hypothetical
syllogism

$$\begin{array}{l} p \vee q \\ \neg p \vee r \\ \hline \therefore q \vee r \end{array}$$

Resolution

$$\begin{array}{l} \neg q \\ p \rightarrow q \\ \hline \therefore \neg p \end{array}$$

Modus
tollens

$$\begin{array}{l} p \vee q \\ \neg p \\ \hline \therefore q \end{array}$$

Disjunctive
syllogism

$$\begin{array}{l} p \\ q \\ \hline \therefore p \wedge q \end{array}$$

Conjunction

Inference Example

- Given:
 - It is humid
 - If it is humid, then it is hot
 - If it is hot and humid, then it is raining
- Show:
 - It is raining

Review

- Today we:
 - Finished looking at equivalences
 - Discovered predicate logic is, and why it's used
 - Learned what quantifiers are, and how to use them
 - Including how to nest them properly
 - Looked at inference and how to do it
 - This is our first major introduction to doing proofs!