



Adversarial Search (aka Games)

Chapter 5

Some material adopted from notes by Charles R. Dyer, U of Wisconsin-Madison

Why study games?

- Interesting, hard problems requiring minimal “initial structure”
- Clear criteria for success
- Study problems involving {hostile, adversarial, competing} agents and uncertainty of interacting with the natural world
- People have used them to assess their intelligence
- Fun, good, easy to understand, PR potential
- Games often define very large search spaces, e.g. chess 35^{100} nodes in search tree, 10^{40} legal states

Chess early days



- **1948:** Norbert Wiener [describes](#) how chess program can work using minimax search with an evaluation function
- **1950:** Claude Shannon publishes [Programming a Computer for Playing Chess](#)
- **1951:** Alan Turing develops *on paper* 1st program capable of playing full chess games ([Turochamp](#))
- **1958:** 1st program plays full game [on IBM 704](#) (loses)
- **1962:** [Kotok & McCarthy](#) (MIT) 1st program to play credibly
- **1967:** Greenblatt's [Mac Hack Six](#) (MIT) defeats a person in regular tournament play

State of the art

- **1979 Backgammon:** [BKG](#) (CMU) tops world champ
- **1994 Checkers:** [Chinook](#) is the world champion
- **1997 Chess:** IBM [Deep Blue](#) beat Gary Kasparov
- **2007 Checkers:** [solved](#) (it's a draw)
- **2016 Go:** [AlphaGo](#) beat champion Lee Sedol
- **2017 Poker:** CMU's [Libratus](#) won \$1.5M from 4 top poker players in 3-week challenge in casino
- **20?? Bridge:** Expert [bridge programs](#) exist, but no world champions yet

AlphaGo - The Movie

Highly recommended 2017 award-winning documentary, **free on YouTube**



**How can
we do it?**

Classical vs. Statistical approach

- We'll look first at the classical approach used from the 1940s to 2010
- Then at newer statistical approaches, of which AlphaGo is an example
- These share some techniques

Typical simple case for a game

- **2-person** game
- Players **alternate moves**
- **Zero-sum**: one player's loss is the other's gain
- **Perfect information**: both players have access to complete information about state of game. No information hidden from either player.
- **No chance** (e.g., using dice) involved
- Examples: Tic-Tac-Toe, Checkers, Chess, Go, Nim, Othello
- But not: Bridge, Solitaire, Backgammon, Poker, Rock-Paper-Scissors, ...

Can we use ...



- Uninformed search?
- Heuristic search?
- Local search?
- Constraint based search?

None of these model the fact
that we have an **adversary** ...

How to play a game

- A way to play such a game is to:
 - Consider all the legal moves you can make
 - Compute new position resulting from each move
 - Evaluate each to determine which is best
 - Make that move
 - Wait for your opponent to move and repeat
- Key problems are:
 - Representing the “board” (i.e., game state)
 - Generating all legal next boards
 - Evaluating a position

Evaluation function

- **Evaluation function** or **static evaluator** used to evaluate the “goodness” of a game position
 - Contrast with heuristic search, where evaluation function estimates **cost** from start node to goal passing through given node
- Zero-sum assumption permits single function to describe goodness of board for both players
 - $f(n) \gg 0$: position n good for me; bad for you
 - $f(n) \ll 0$: position n bad for me; good for you
 - $f(n)$ near 0 : position n is a neutral position
 - $f(n) = +\text{infinity}$: win for me
 - $f(n) = -\text{infinity}$: win for you

Evaluation function examples

- **For Tic-Tac-Toe**

$$f(n) = [\# \text{ my open 3lengths}] - [\# \text{ your open 3lengths}]$$

Where a 3length is complete row, column or diagonal that has no opponent marks

- **Alan Turing's function for chess**

- $f(n) = w(n)/b(n)$ where $w(n)$ = sum of point value of white's pieces and $b(n)$ = sum of black's

- Traditional piece values: pawn:1; knight:3; bishop:3; rook:5; queen:9

Evaluation function examples

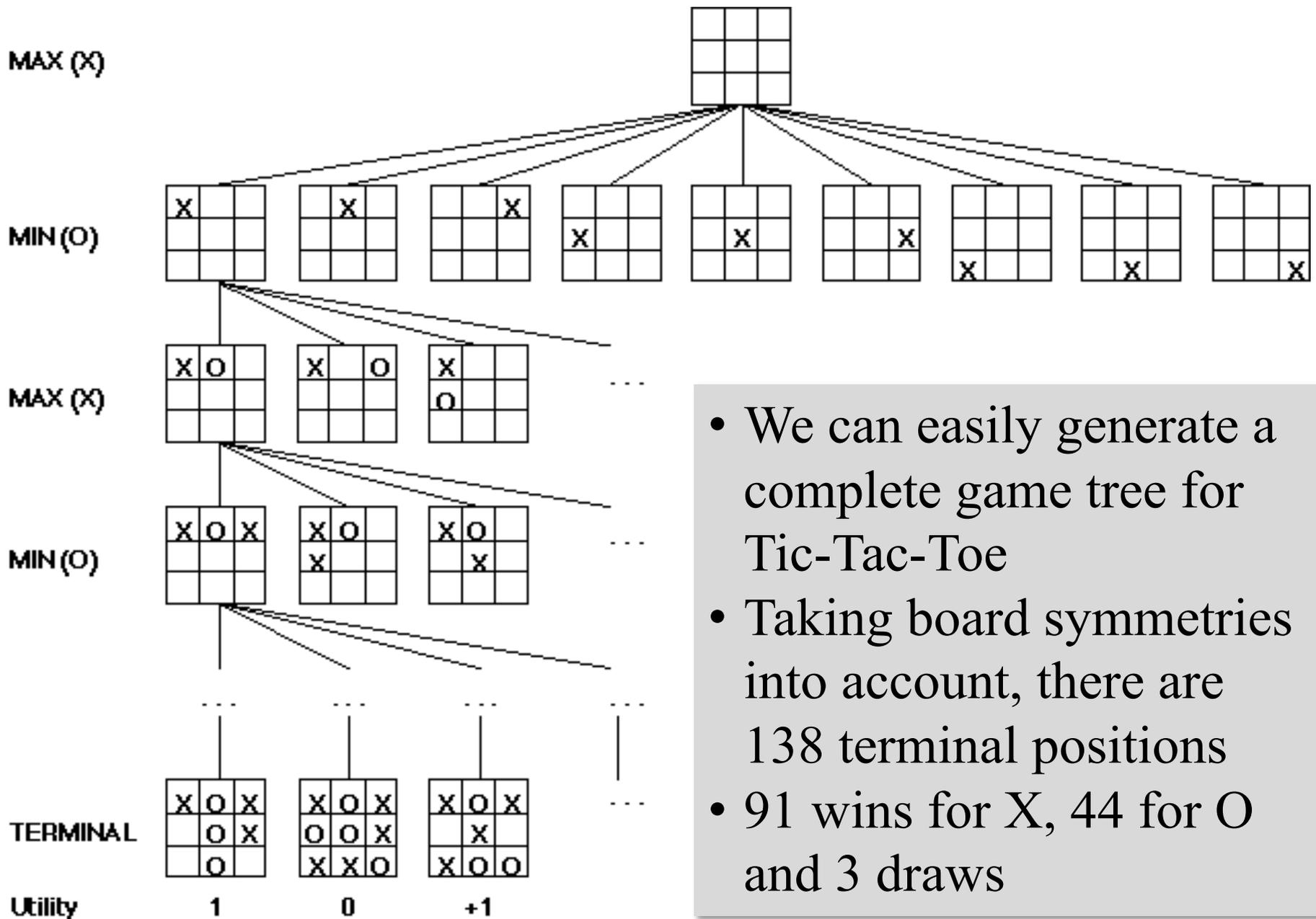
- Most evaluation functions specified as a weighted sum of positive features

$$f(n) = w_1 * \text{feat}_1(n) + w_2 * \text{feat}_2(n) + \dots + w_n * \text{feat}_k(n)$$

- Typical chess features are piece count, piece values, piece placement, squares controlled, etc.
- IBM's chess program [Deep Blue](#) (circa 1996) had >8K features in its evaluation function

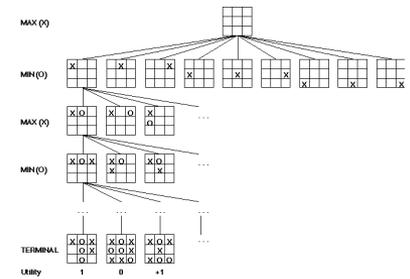
But, that's not how people play

- People also use *look ahead*
i.e., enumerate actions, consider opponent's possible responses, REPEAT
- Producing a *complete* game tree is only possible for simple games
- So, generate a partial game tree for some number of plys
 - Move = each player takes a turn
 - Ply = one player's turn
- What do we do with the game tree?



- We can easily generate a complete game tree for Tic-Tac-Toe
- Taking board symmetries into account, there are 138 terminal positions
- 91 wins for X, 44 for O and 3 draws

Game trees



- Problem spaces for typical games are trees
- Root node is current board configuration; player must decide best single move to make next
- **Static evaluator function** rates board position **f(board):real**, often >0 for me; <0 for opponent
- Arcs represent possible legal moves for a player
- If **my turn** to move, then root is labeled a "**MAX**" node; otherwise it's a "**MIN**" node
- Each tree level's nodes are all MAX or all MIN; nodes at level i are of opposite kind from those at level $i+1$

Minimax procedure

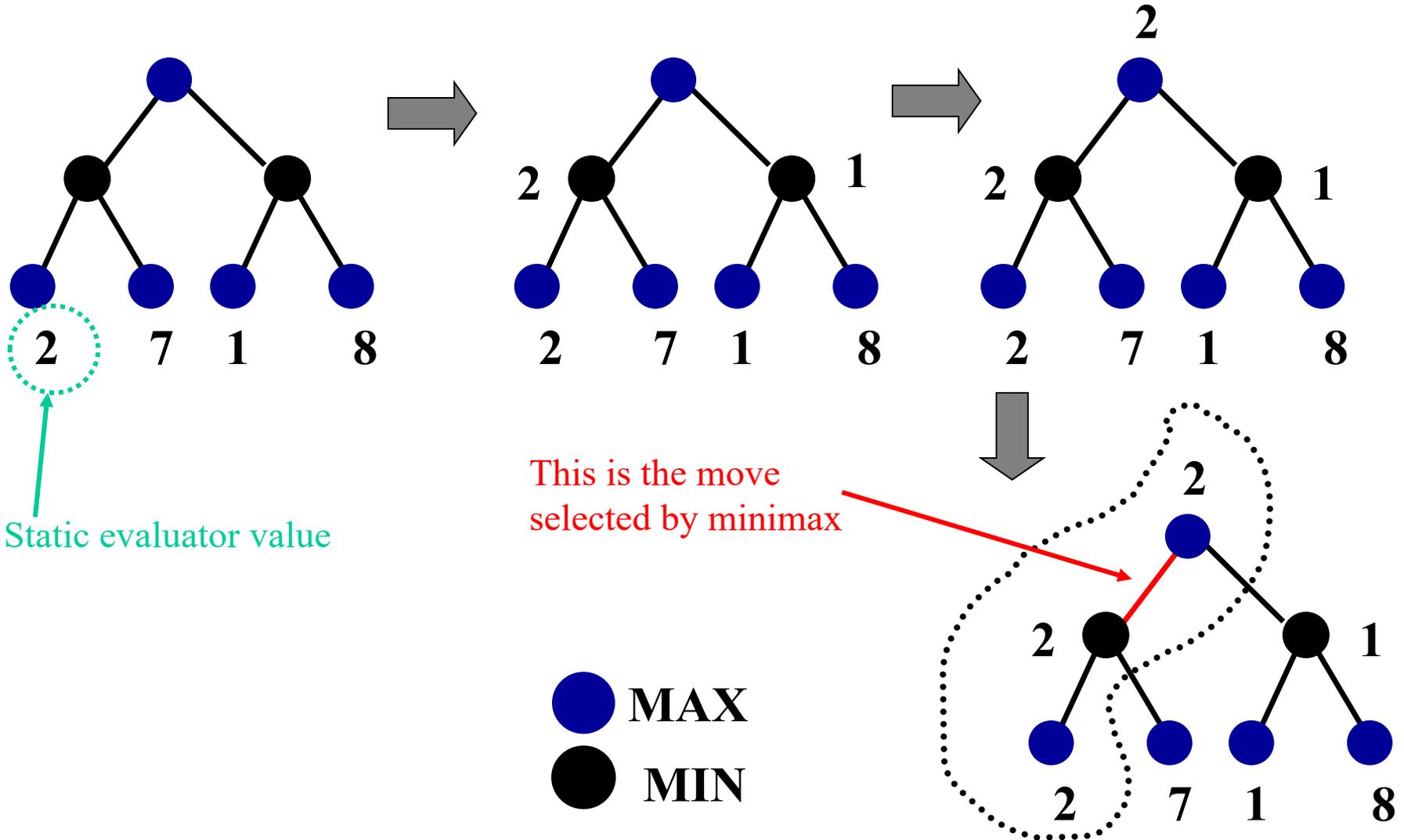
- Create MAX node with current board configuration
- Expand nodes to some **depth** (a.k.a. **plys**) of lookahead in game
- Apply evaluation function at each **leaf** node
- *Back up* values for each non-leaf node until value is computed for the root node
 - At MIN nodes: value is **minimum** of children's values
 - At MAX nodes: value is **maximum** of children's values
- Choose move to child node whose backed-up value determined value at root

Minimax theorem

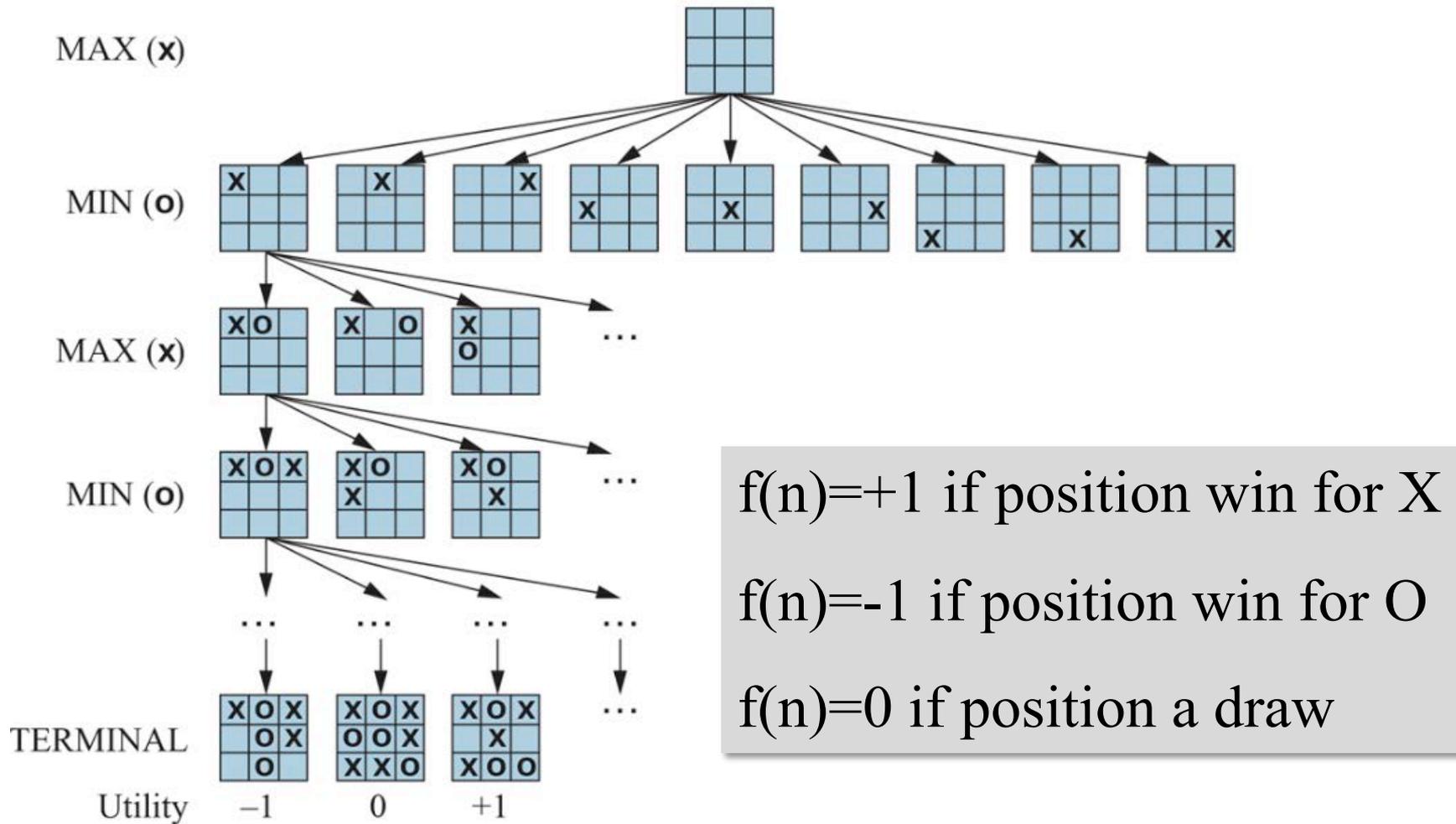
- Intuition: assume your opponent is at least as smart as you and play accordingly
 - If she's not, you can only do better!
- [Von Neumann](#), J: *Zur Theorie der Gesellschaftsspiele* Math. Annalen. **100** (1928) 295-320

For every 2-person, 0-sum game with finite strategies, there is a value V and a mixed strategy for each player, such that (a) given player 2's strategy, best payoff possible for player 1 is V , and (b) given player 1's strategy, best payoff possible for player 2 is $-V$.
- You can think of this as:
 - Minimizing your maximum possible loss
 - Maximizing your minimum possible gain

Minimax Algorithm



Partial Game Tree for Tic-Tac-Toe

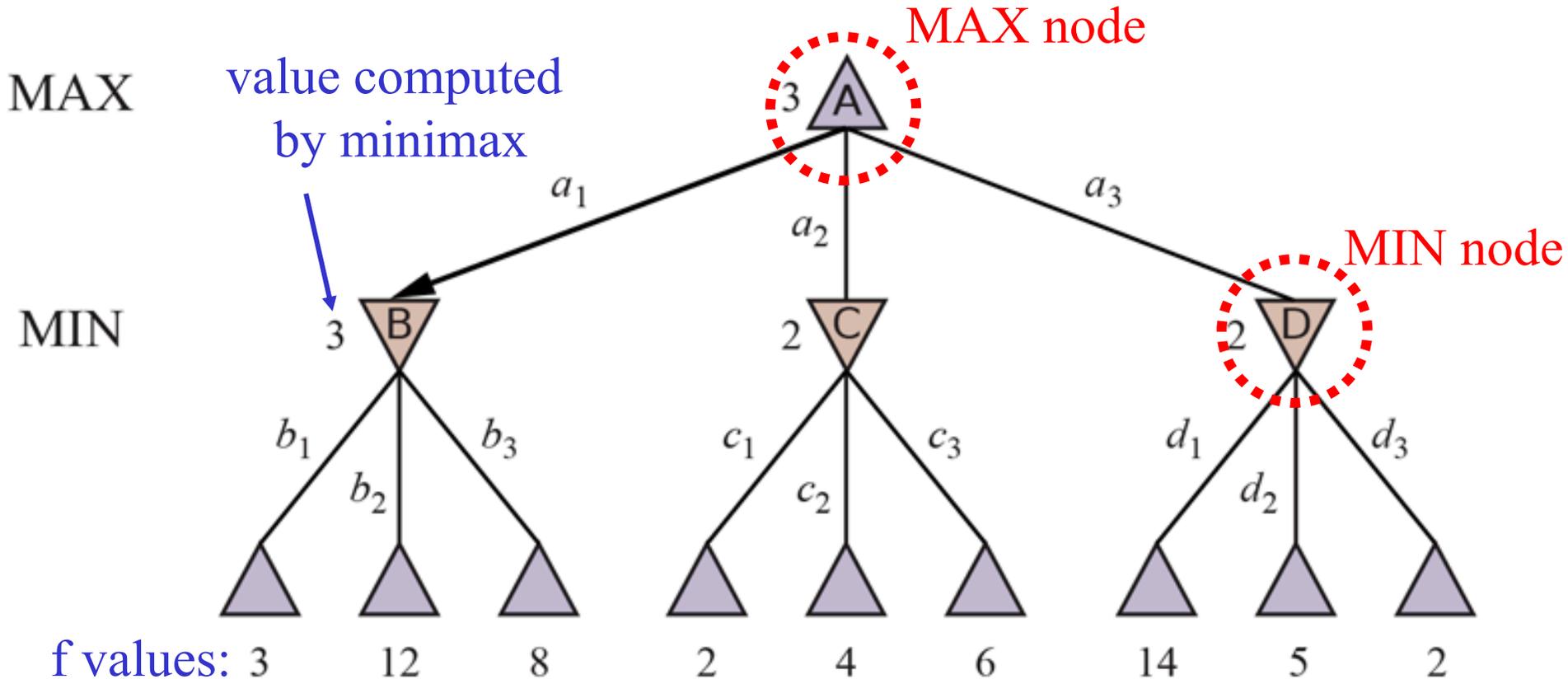


Partial game tree for tic-tac-toe. Top node is the initial state, and max moves first, placing an X in an empty square. Only part of the tree shown, giving alternating moves by min (O) and max (X), until we reach terminal states, which are assigned utilities $\{-1, 0, +1\}$ for $\{\text{lose, draw, win}\}$

Why backed-up values?

- Why not just use a good static evaluator metric on immediate children
- **Intuition:** if metric is good, doing look ahead and backing up values with Minimax should be better
- Non-leaf node N 's backed-up value is value of best state MAX can reach at depth h if MIN plays *well*
 - “plays well”: same criterion as MAX applies to itself
- If e is good, then backed-up value is better estimate of $STATE(N)$ goodness than $e(STATE(N))$
- Use lookahead horizon h because time to choose move is limited

Minimax Tree

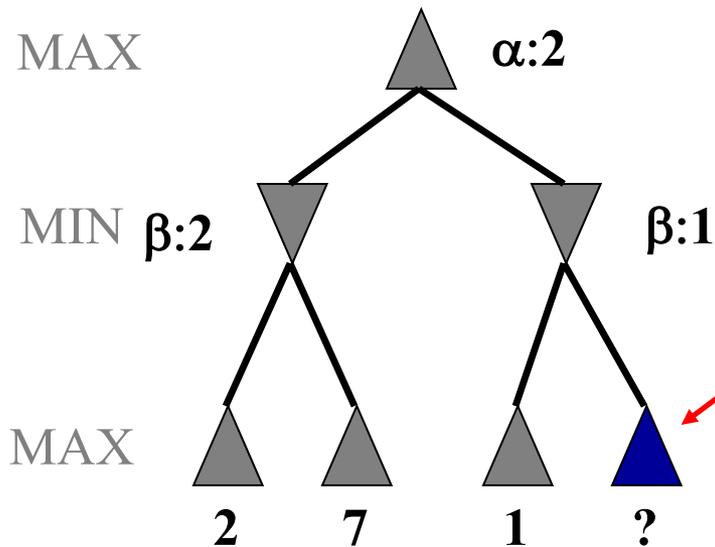


Two-ply game tree. \triangle nodes are “max nodes,” in which it is max’s turn to move, and ∇ nodes are “min nodes.” The terminal nodes show utility values for max; the other nodes are labeled with their minimax values. max’s best move at root is a_1 because it leads to state with the highest minimax value. min’s best reply is b_1 since it leads to the state with the lowest minimax value.

**Is that all
there is to simple
games?**

Alpha-beta pruning

- Improve performance of the minimax algorithm through alpha-beta pruning
- *“If you have an idea that is surely bad, don't take the time to see how truly awful it is”* -Pat Winston (MIT)



- We don't need to compute the value at this node
- No matter what it is, it can't affect value of the root node

Alpha-beta pruning

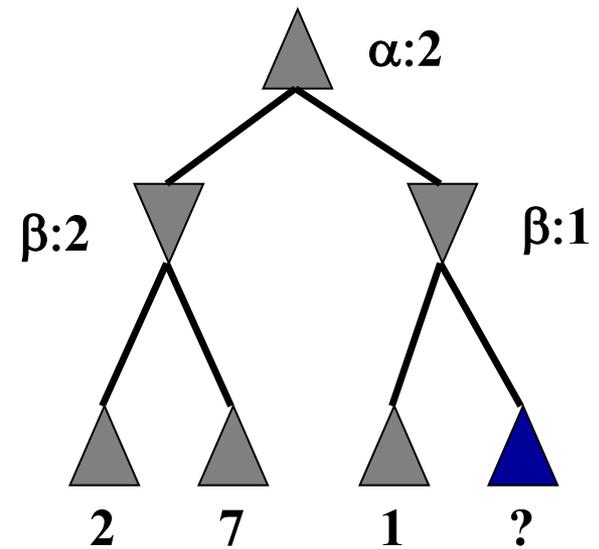
- Traverse tree in depth-first order
- At **MAX** node n , **alpha(n)** = max value found so far

Alpha values start at $-\infty$ and only increase

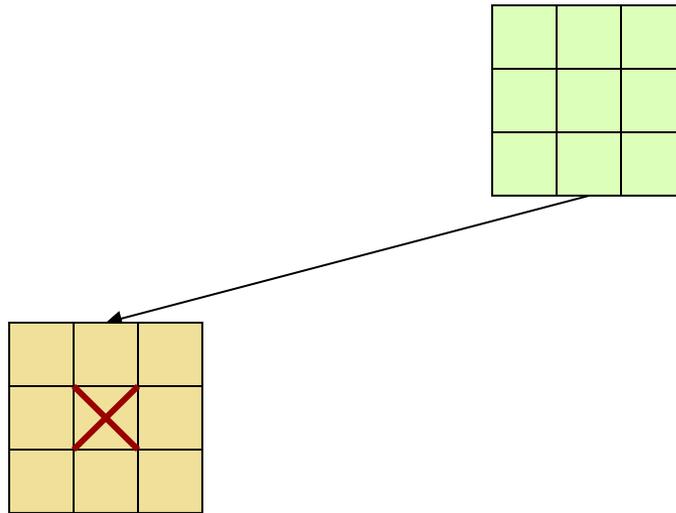
- At **MIN** node n , **beta(n)** = min value found so far

Beta values start at $+\infty$ and only decrease

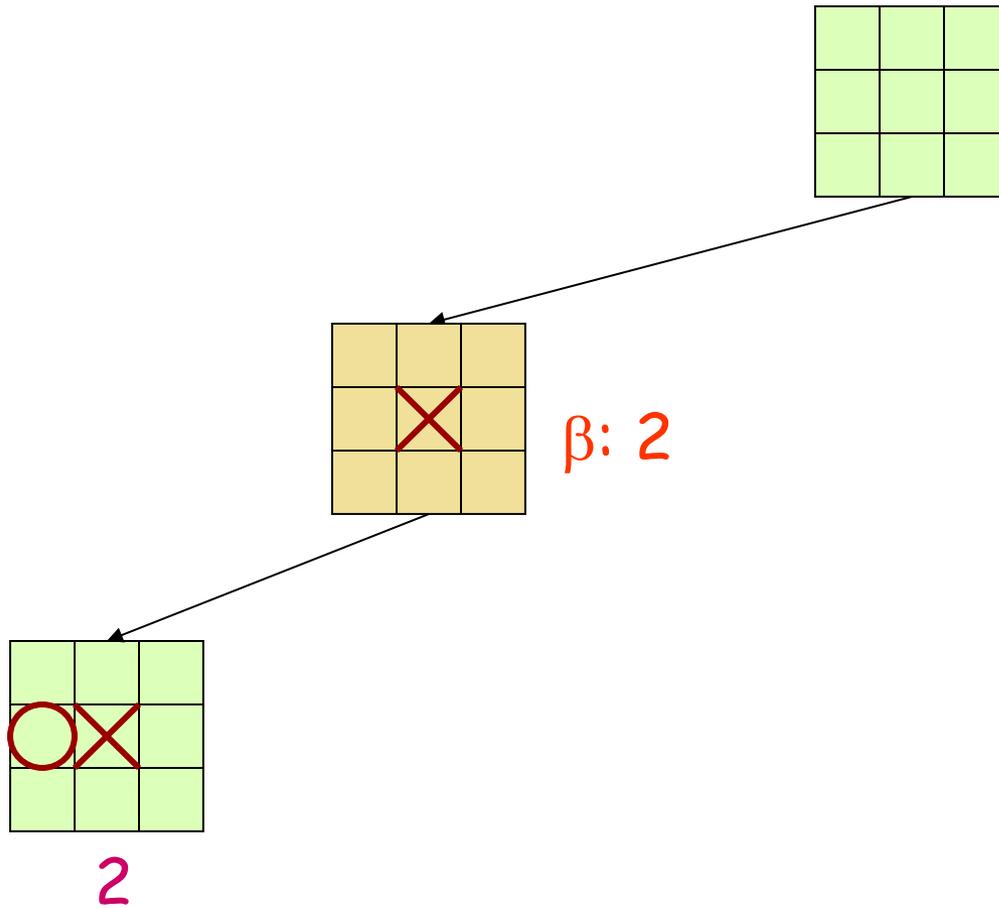
- **Beta cutoff:** stop search below MAX node N (i.e., don't examine more descendants) if $\text{alpha}(N) \geq \text{beta}(i)$ for some MIN node ancestor i of N
- **Alpha cutoff:** stop search below MIN node N if $\text{beta}(N) \leq \text{alpha}(i)$ for a MAX node ancestor i of N



Alpha-Beta Tic-Tac-Toe Example



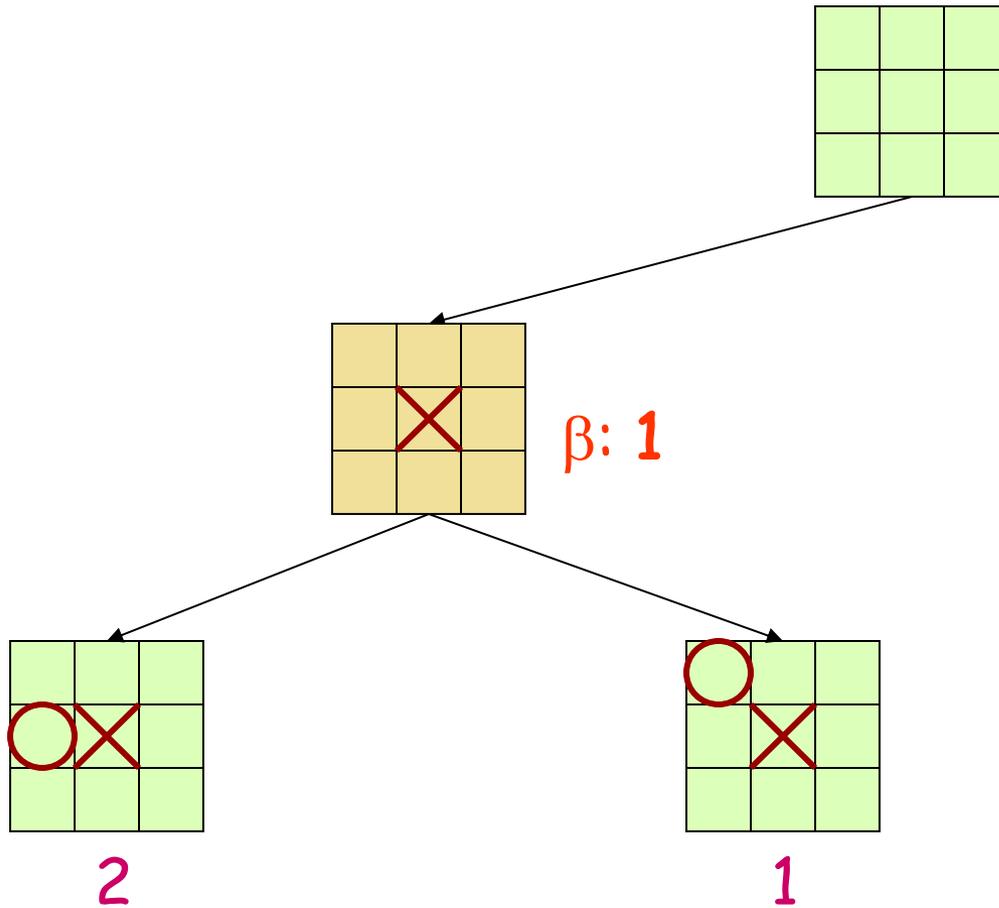
Alpha-Beta Tic-Tac-Toe Example



$F = X$'s open lines –
O's open lines

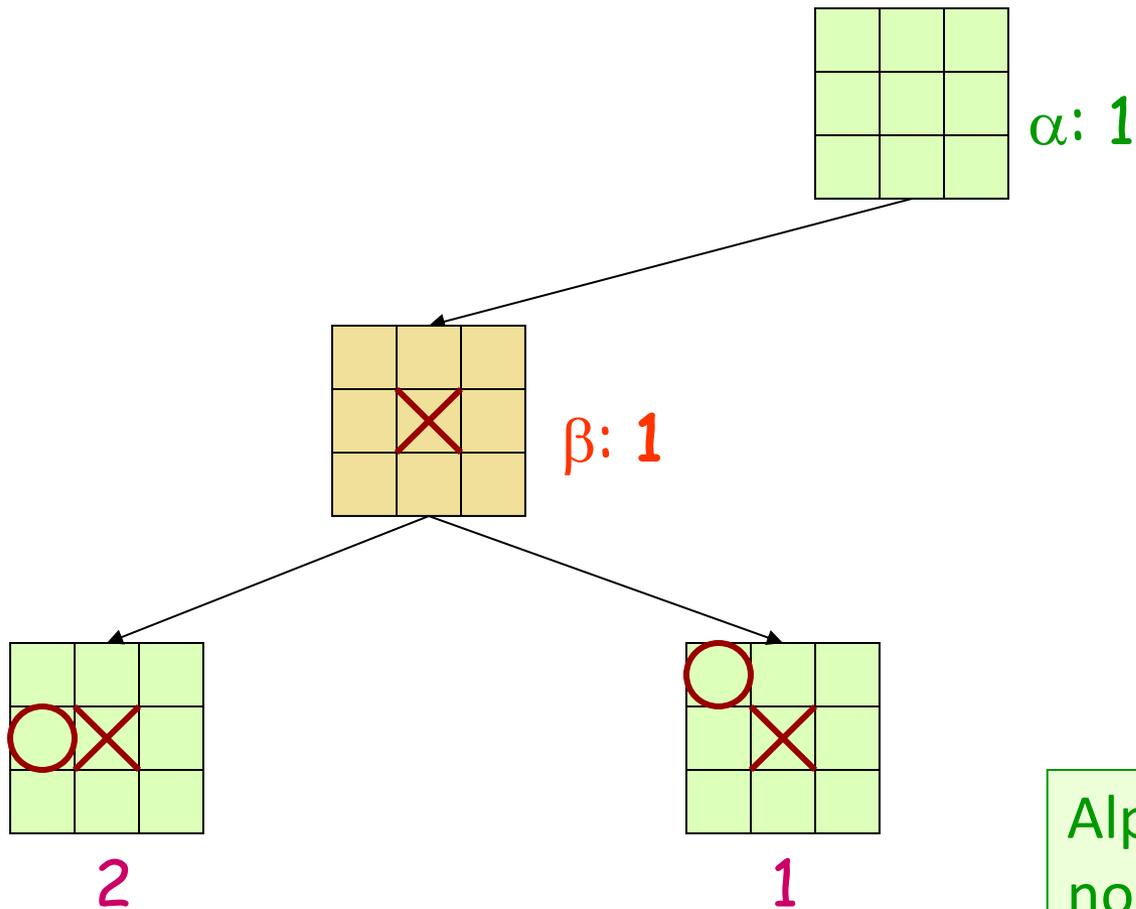
Beta value of a MIN node is **upper** bound on final backed-up value; it can never increase

Alpha-Beta Tic-Tac-Toe Example



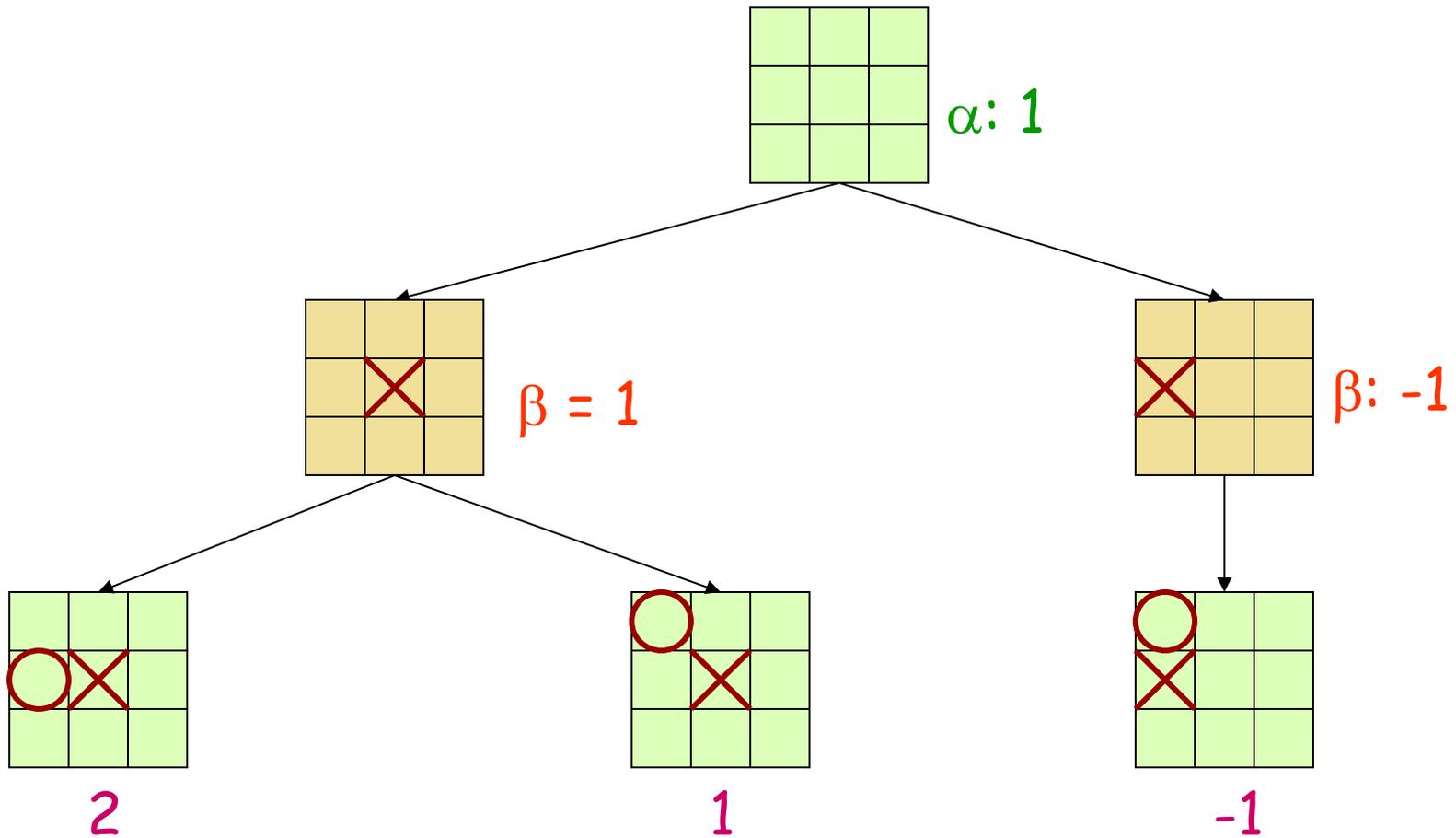
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Alpha-Beta Tic-Tac-Toe Example

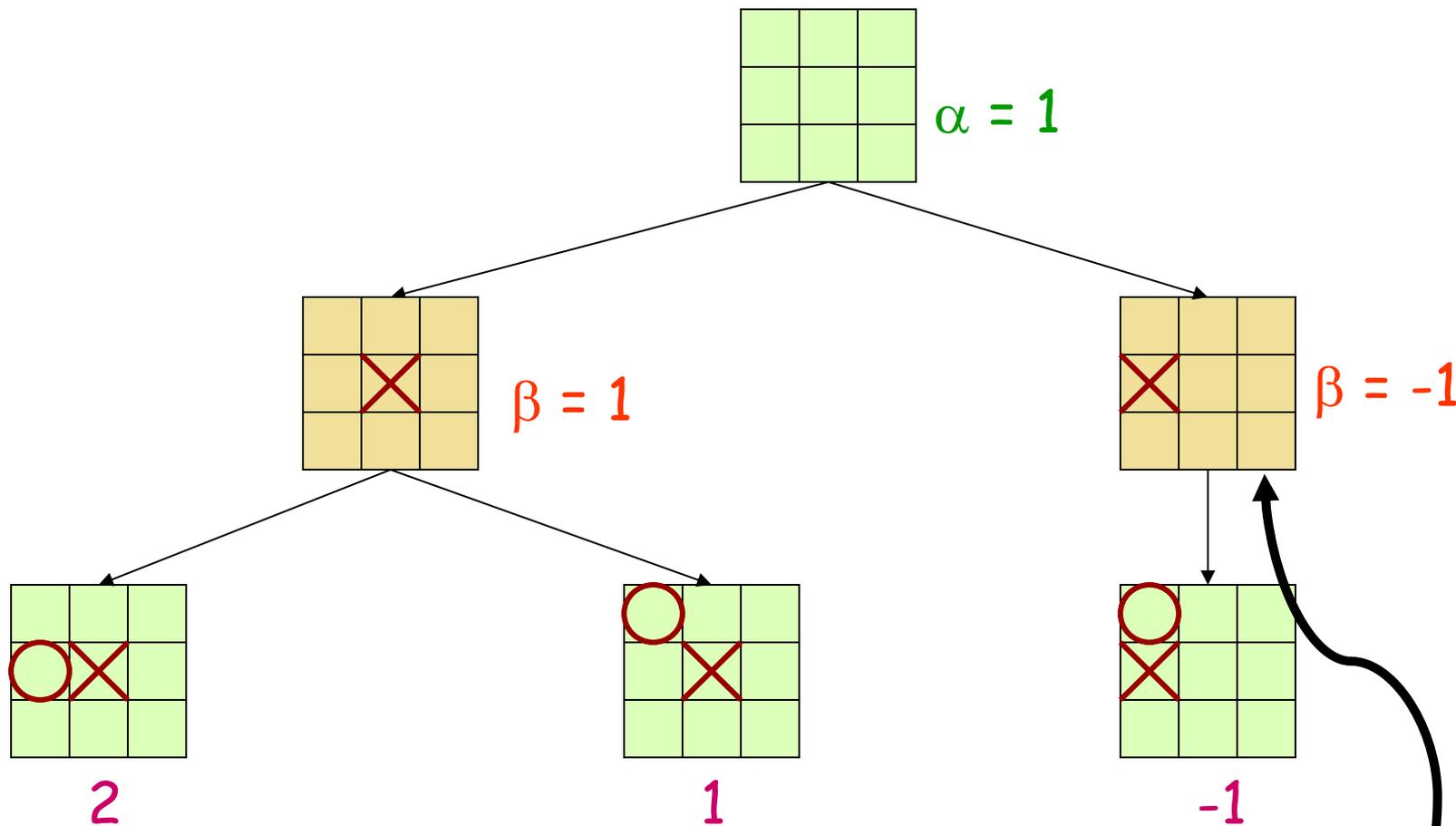


Alpha value of MAX node is **lower** bound on final backed-up value; it can never decrease

Alpha-Beta Tic-Tac-Toe Example

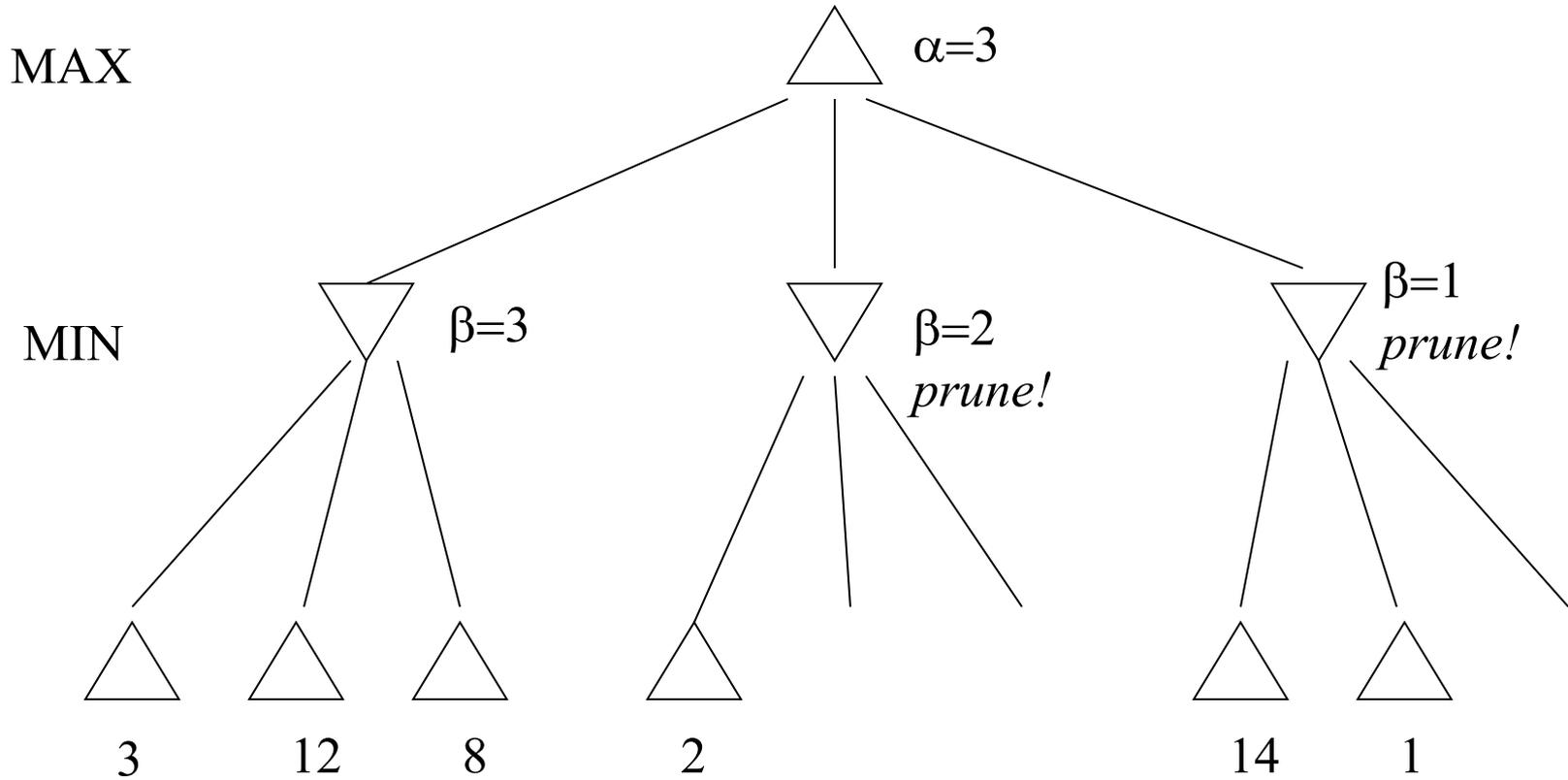


Alpha-Beta Tic-Tac-Toe Example

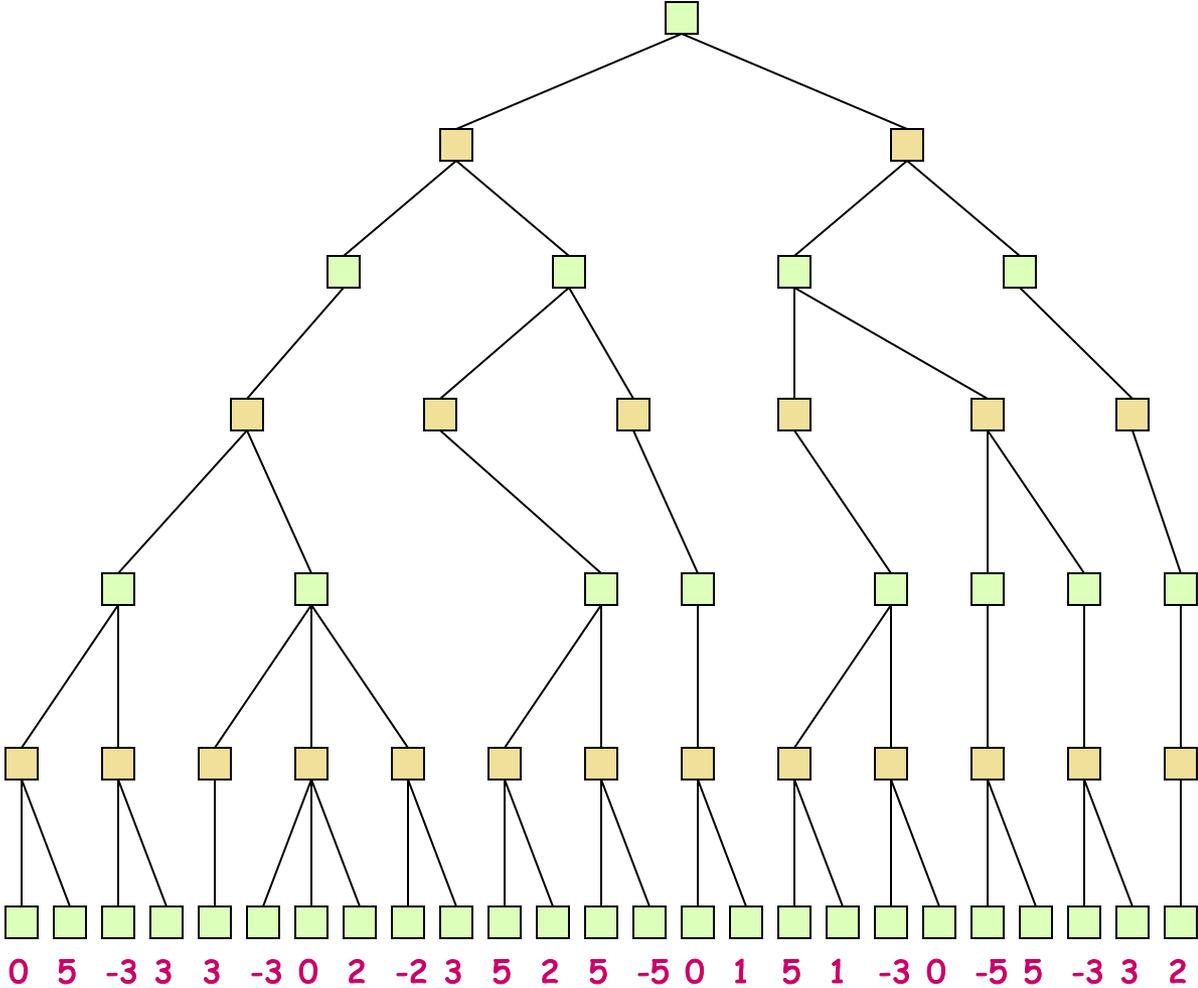


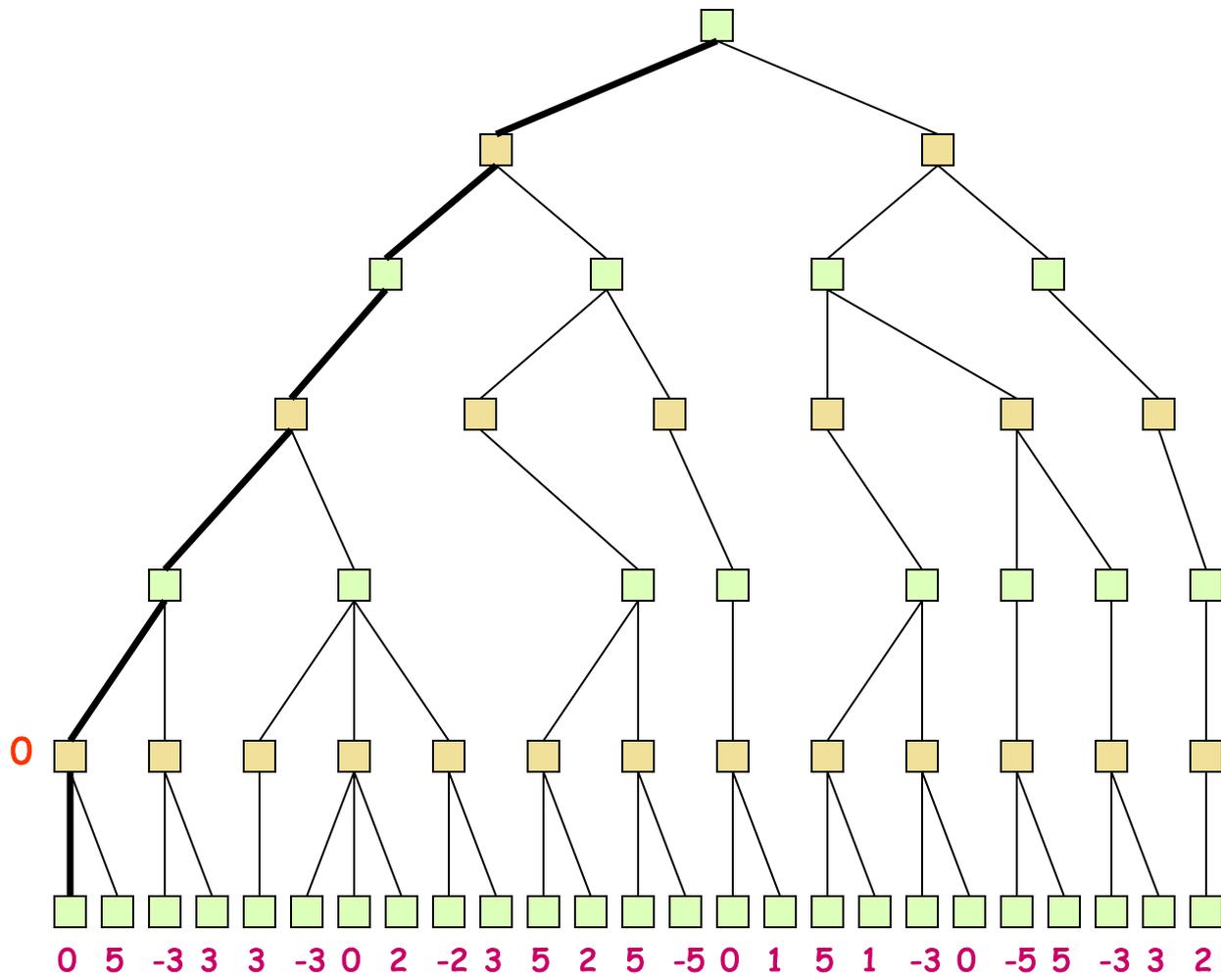
Discontinue search below a MIN node whose beta value \leq alpha value of one of its MAX ancestors

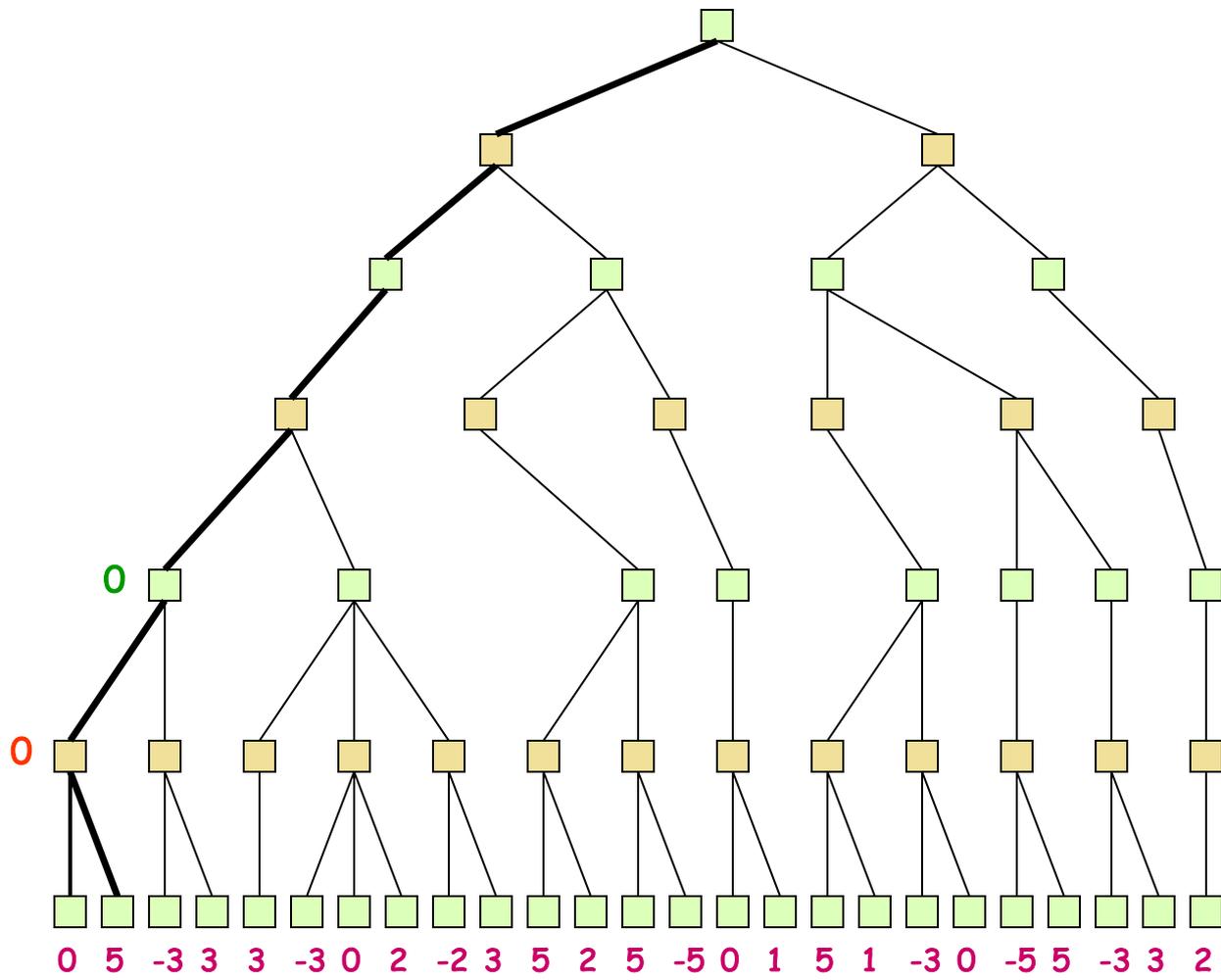
Another alpha-beta example

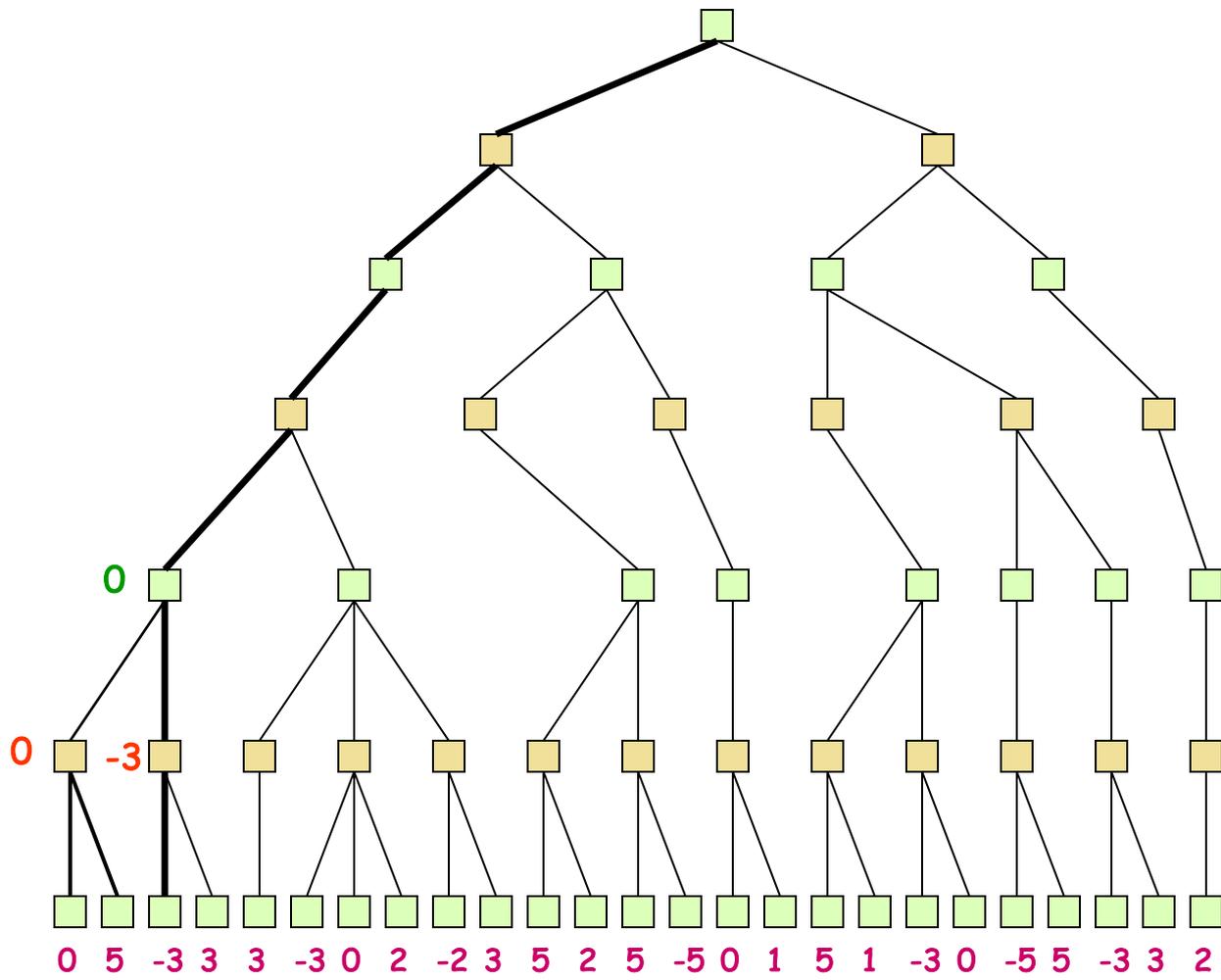


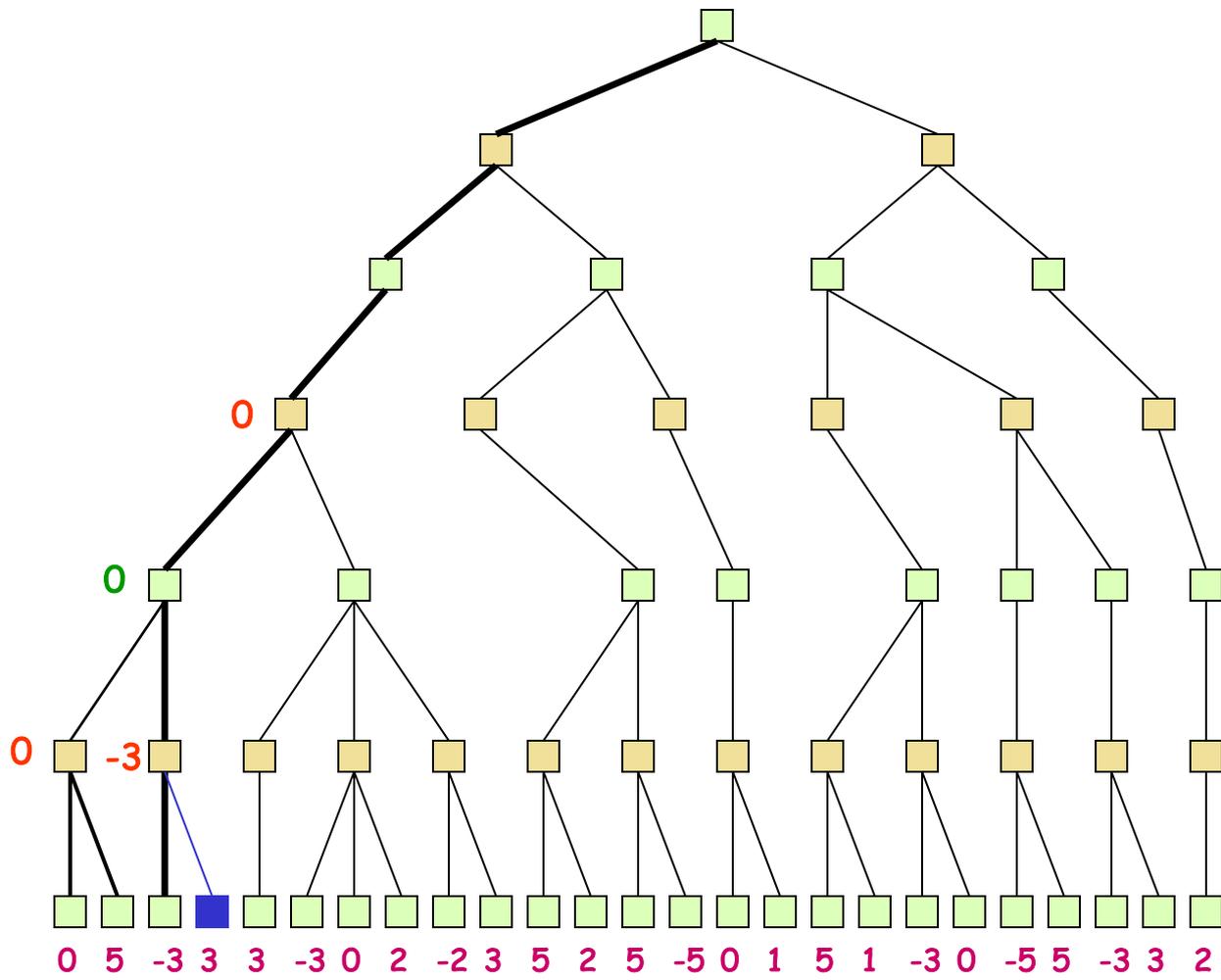
Alpha-Beta Tic-Tac-Toe Example 2

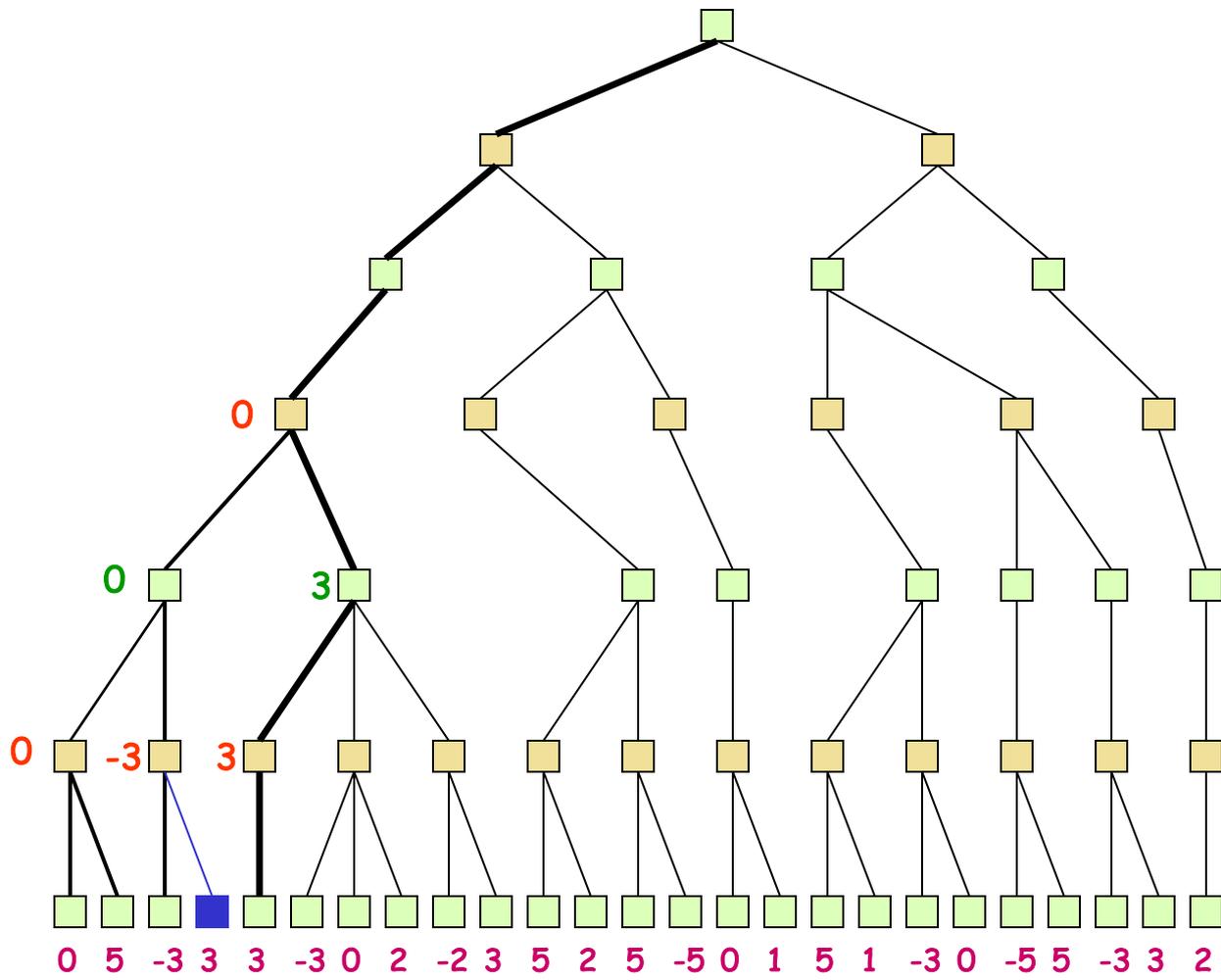


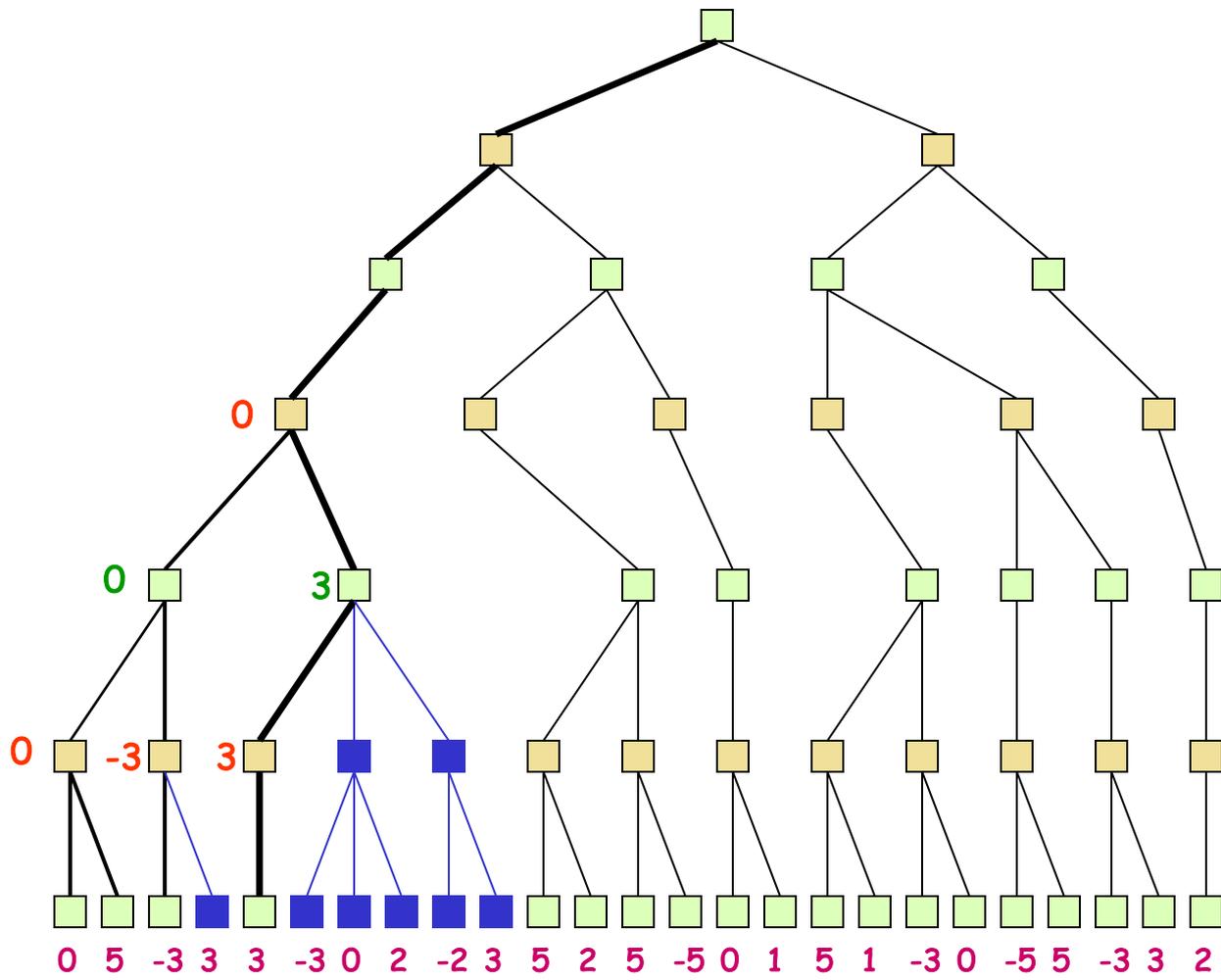


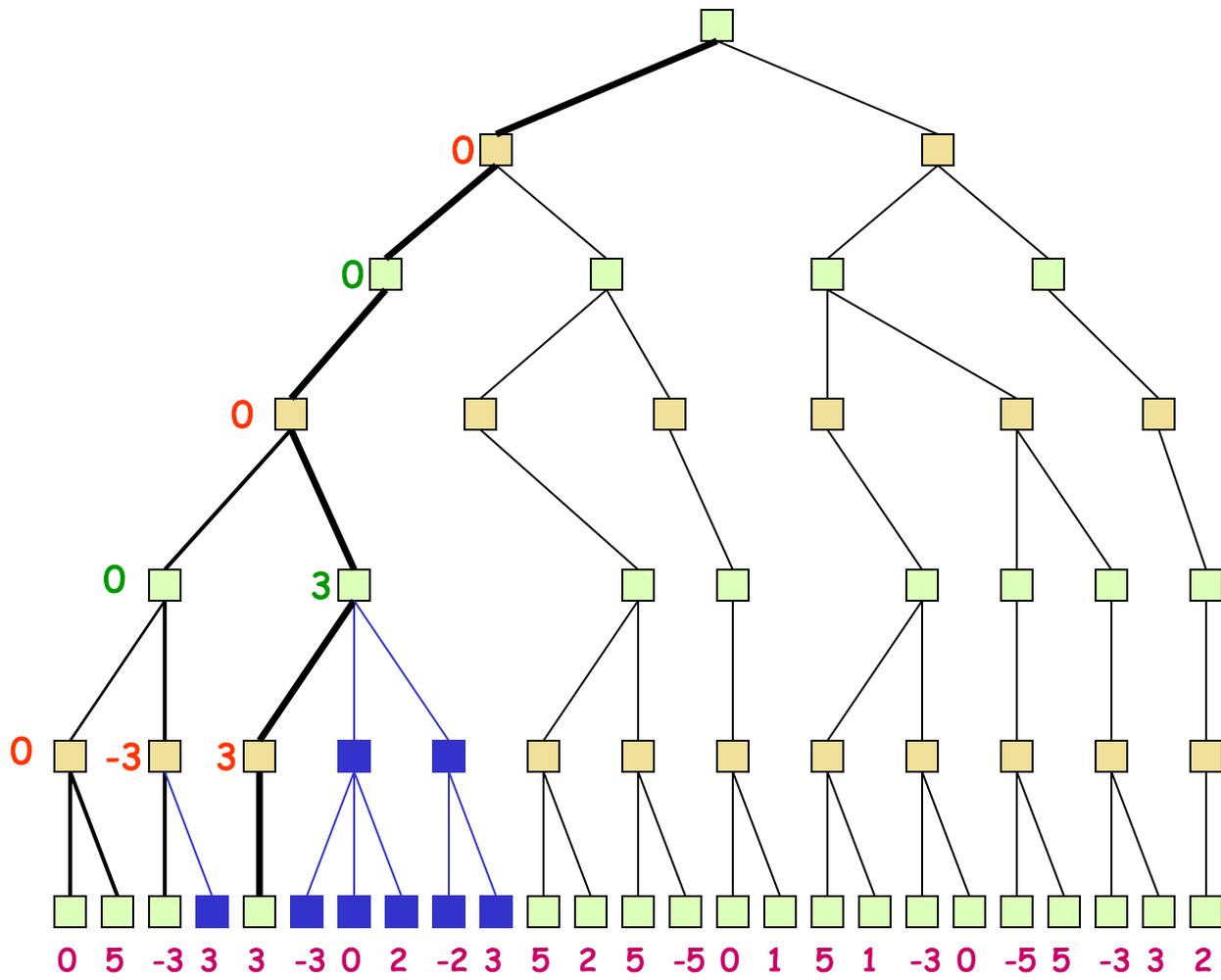


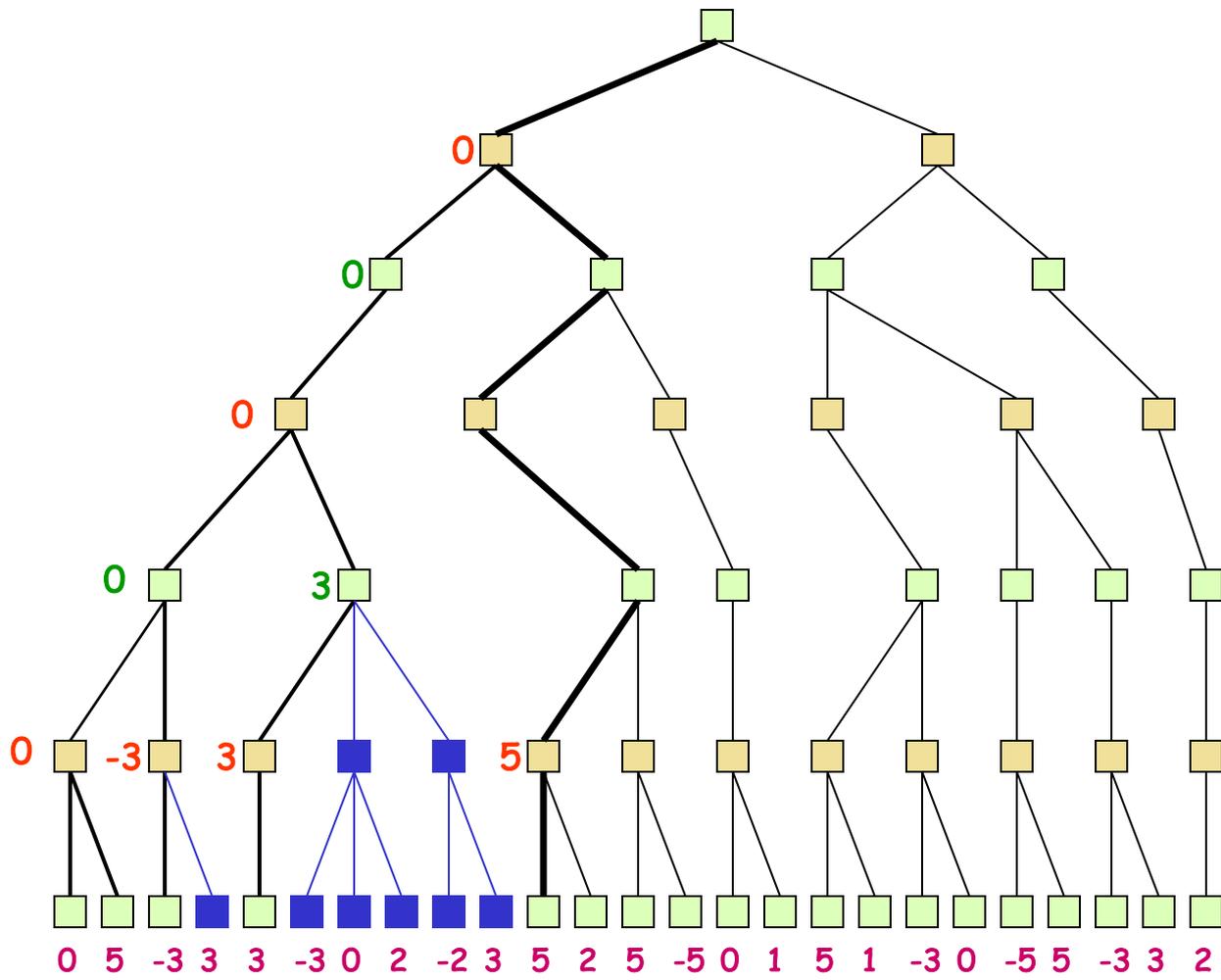


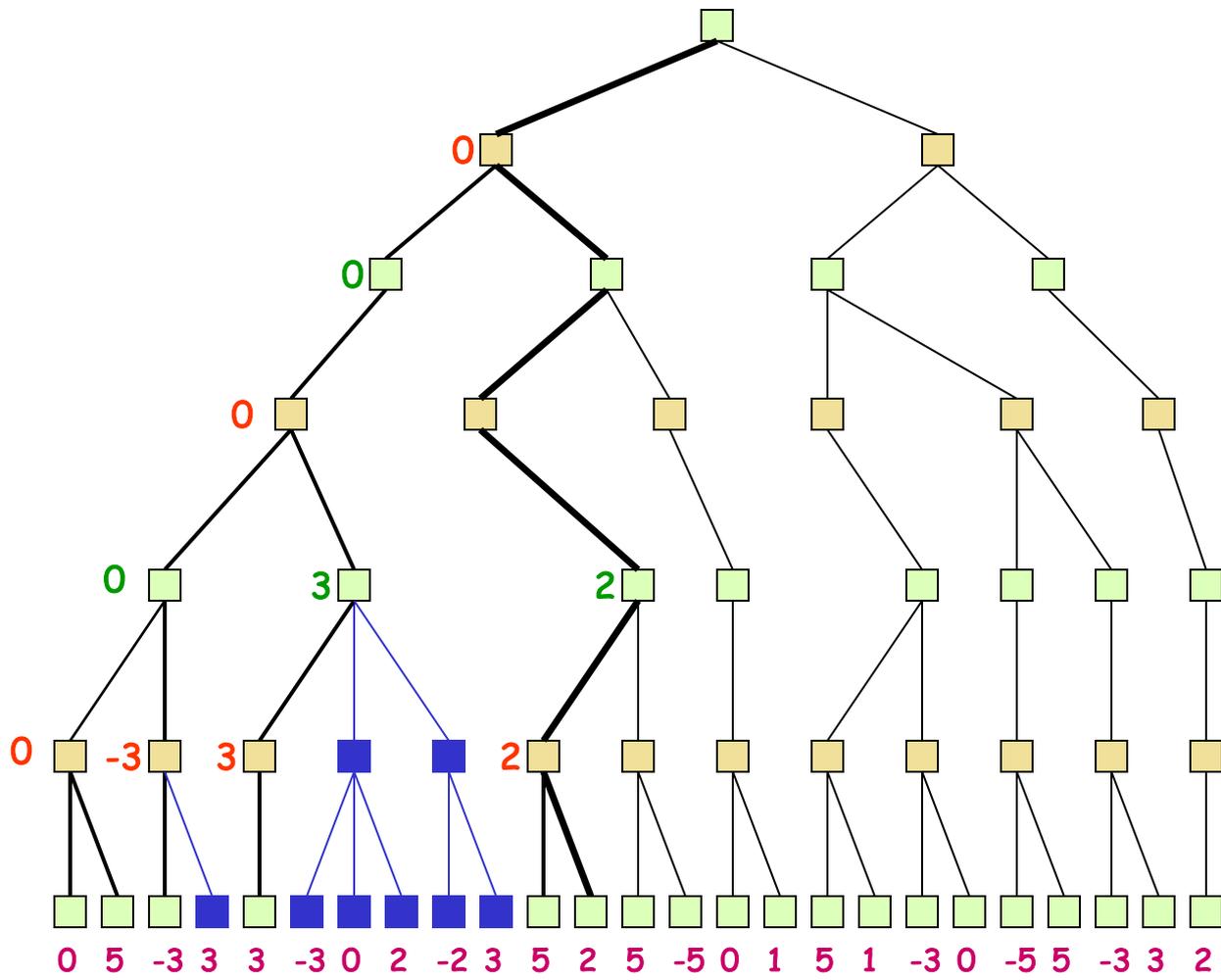


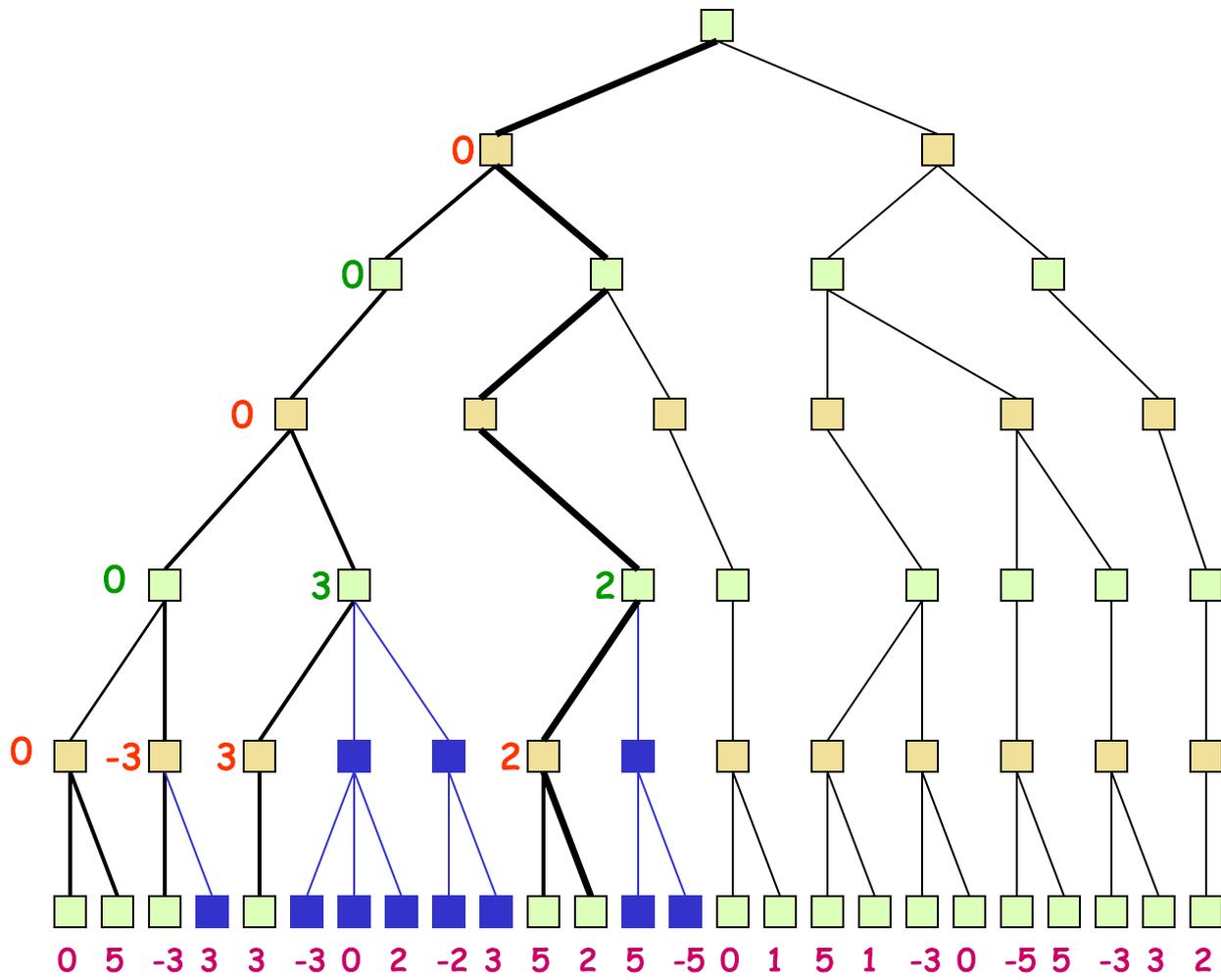


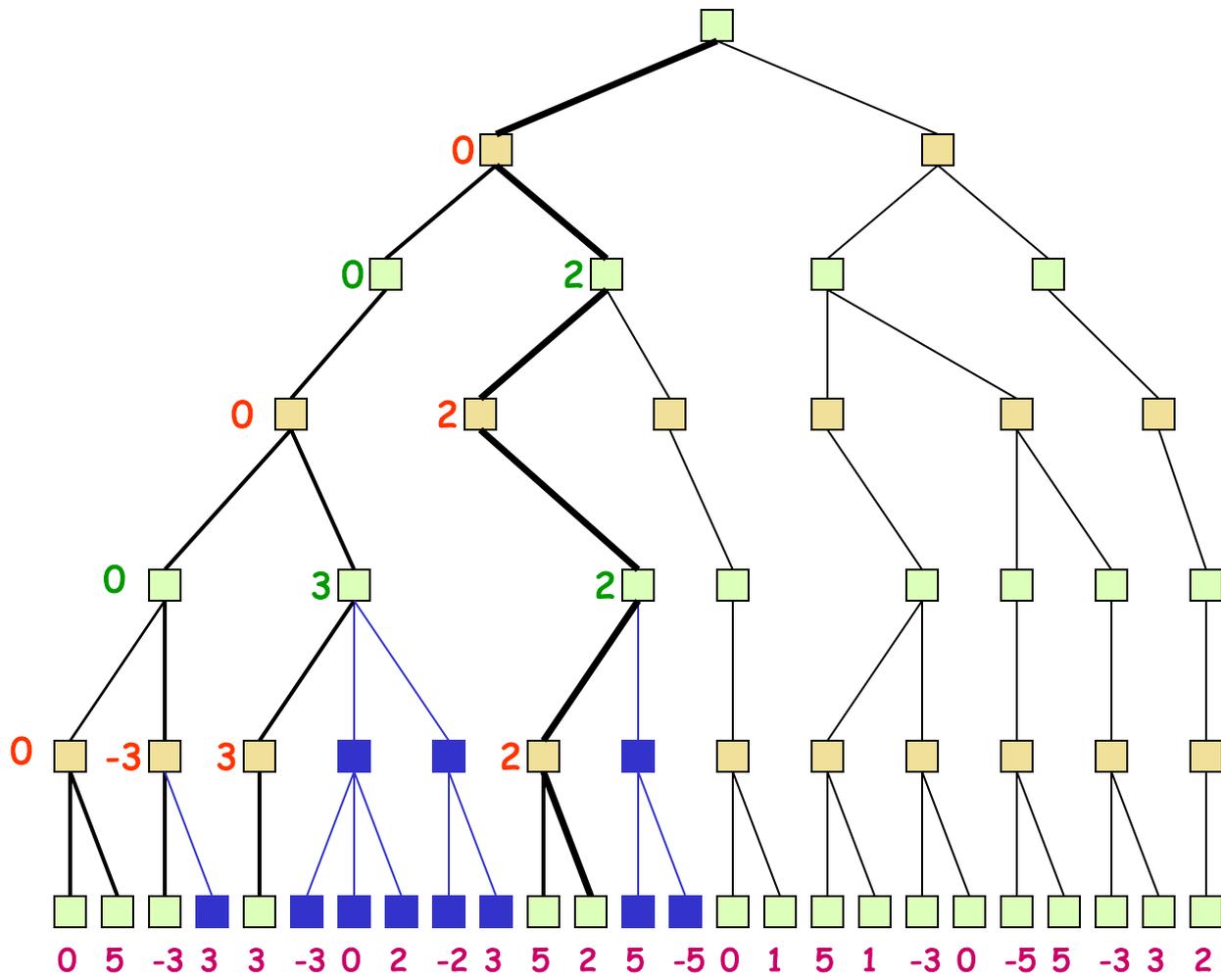


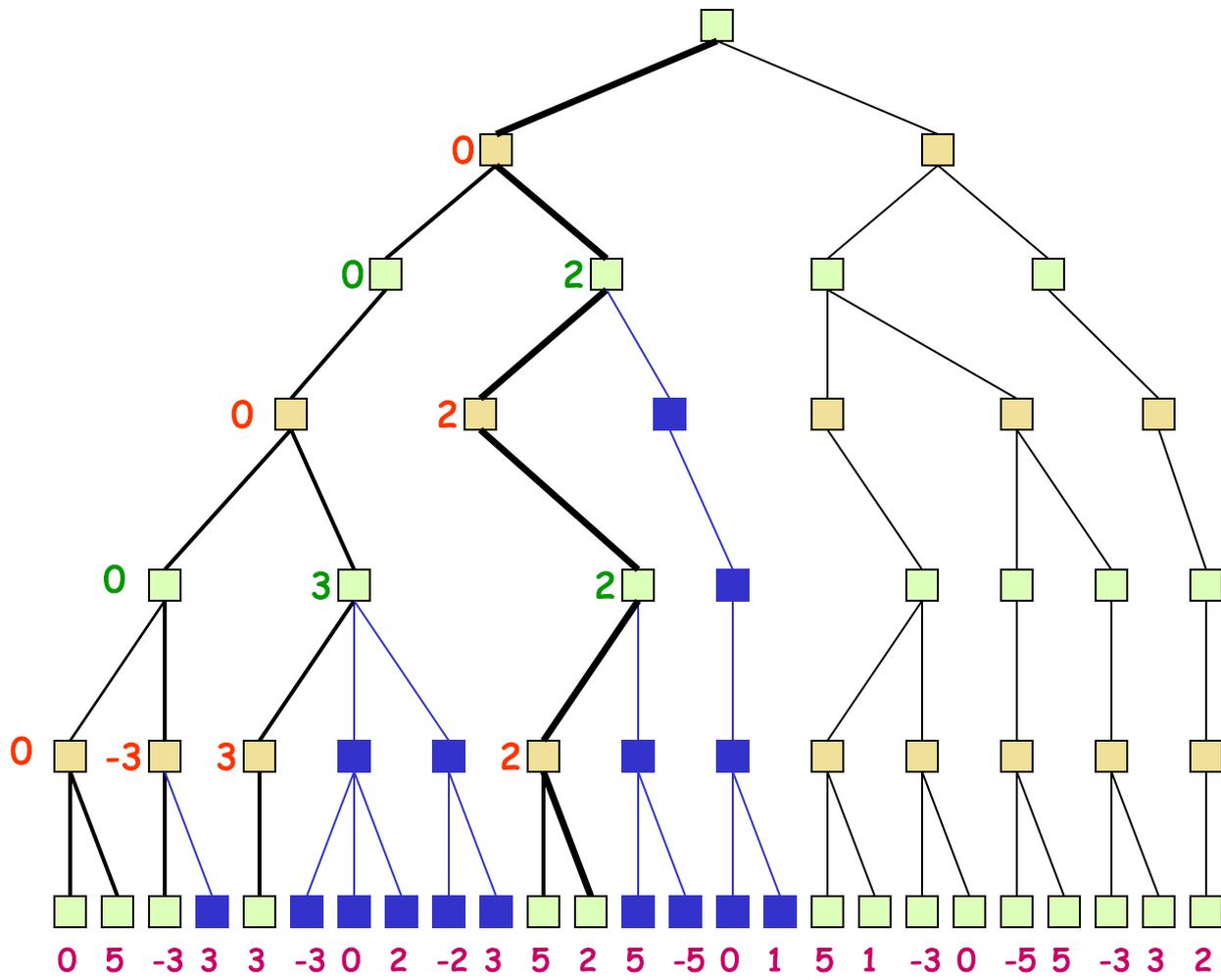


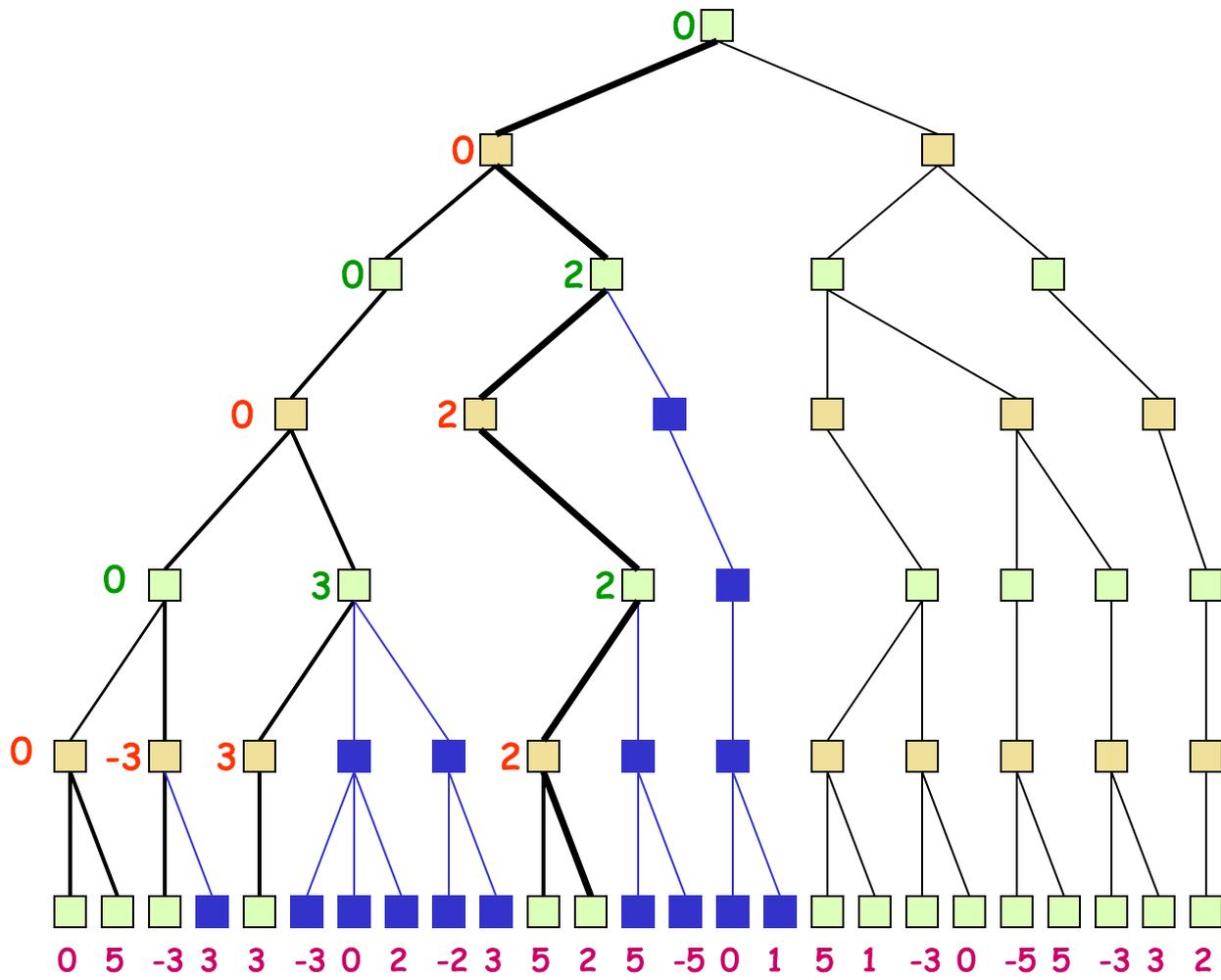


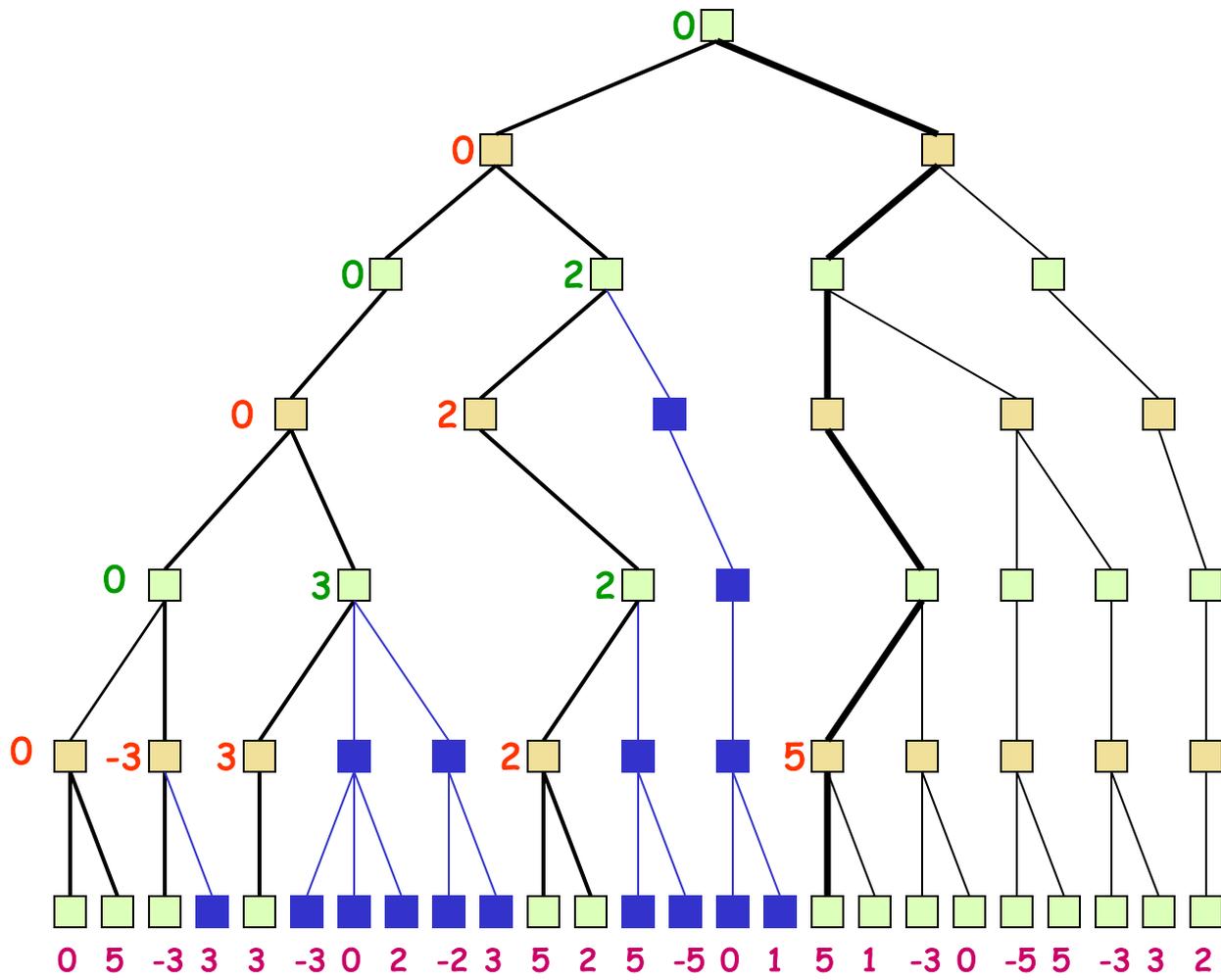


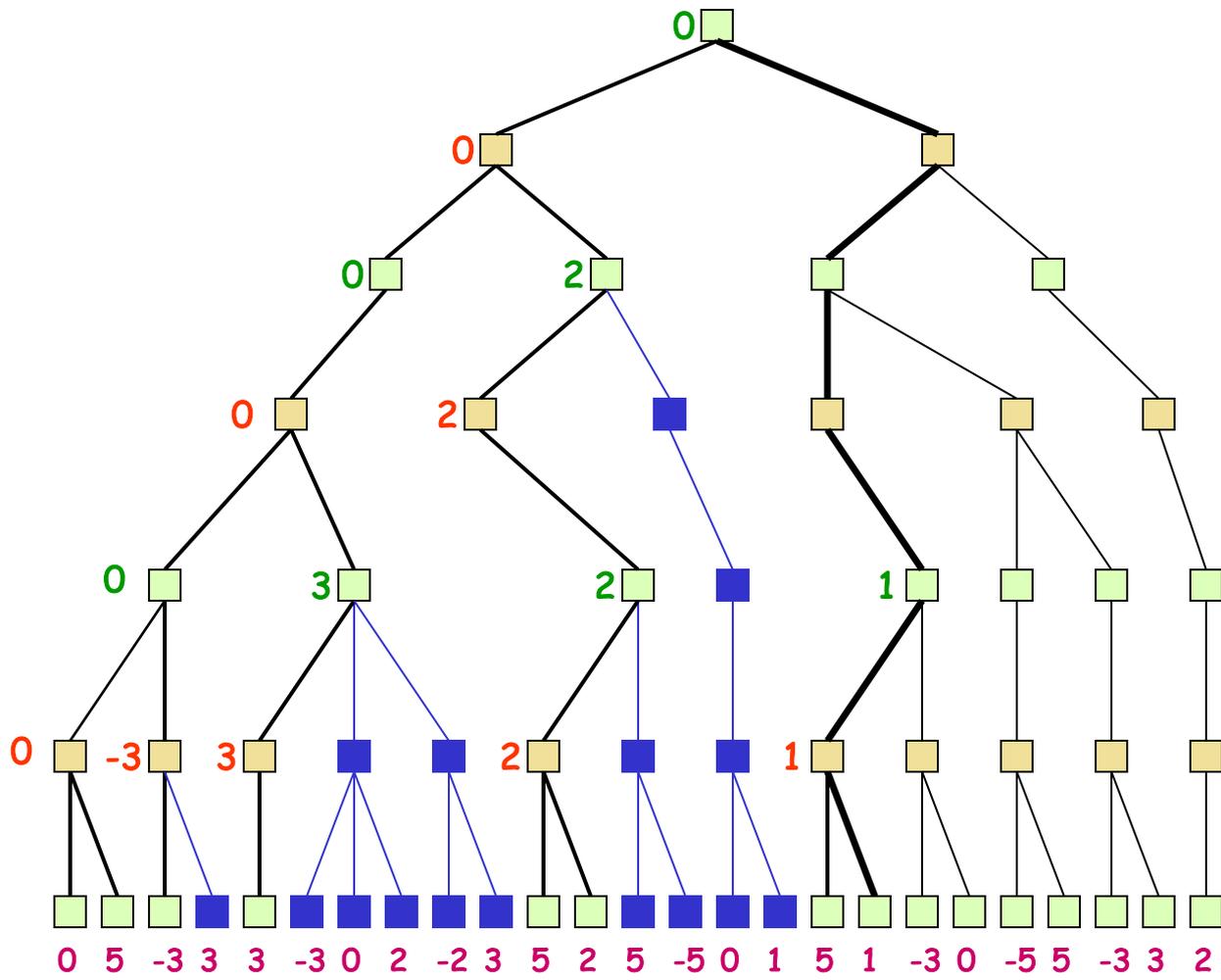


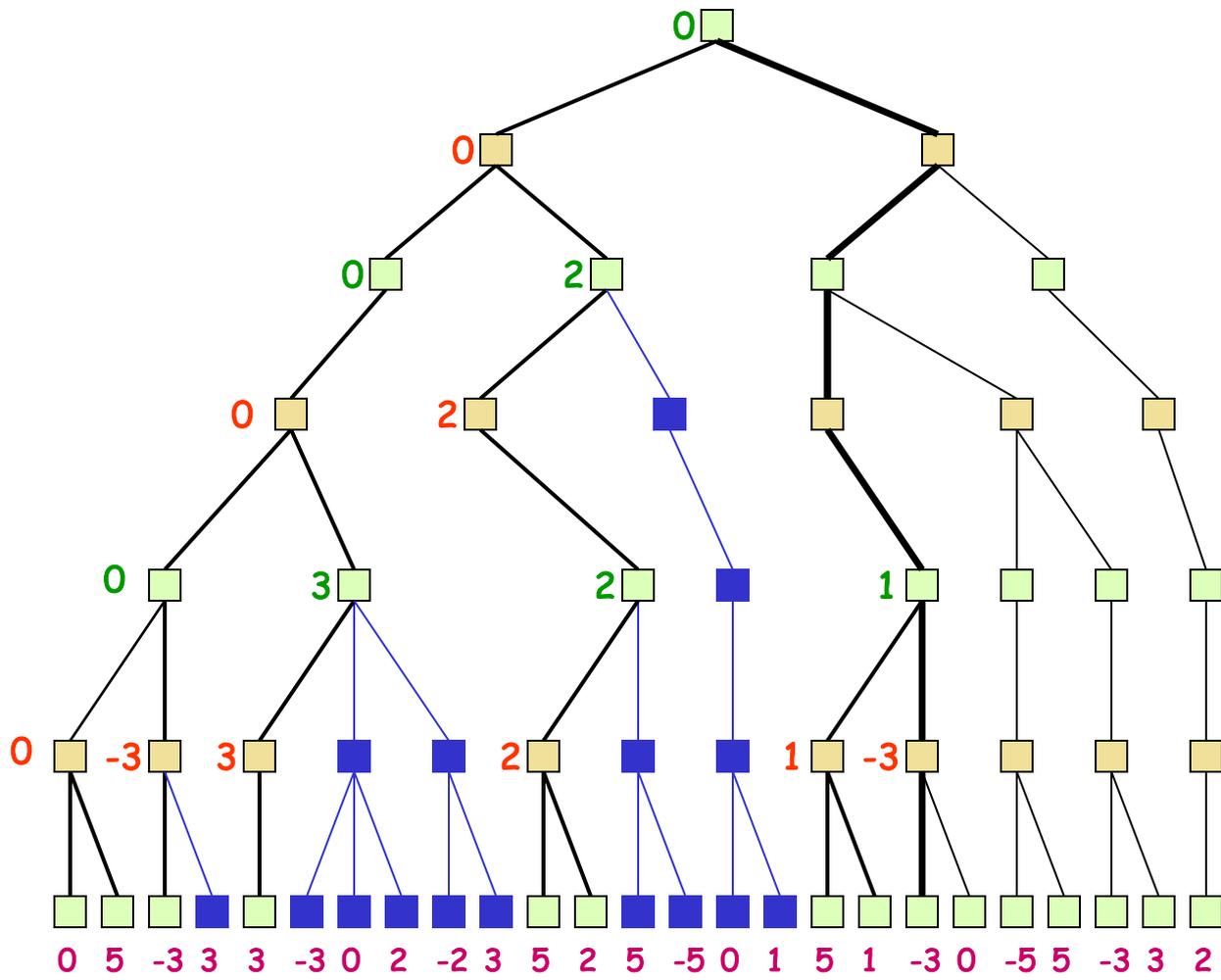


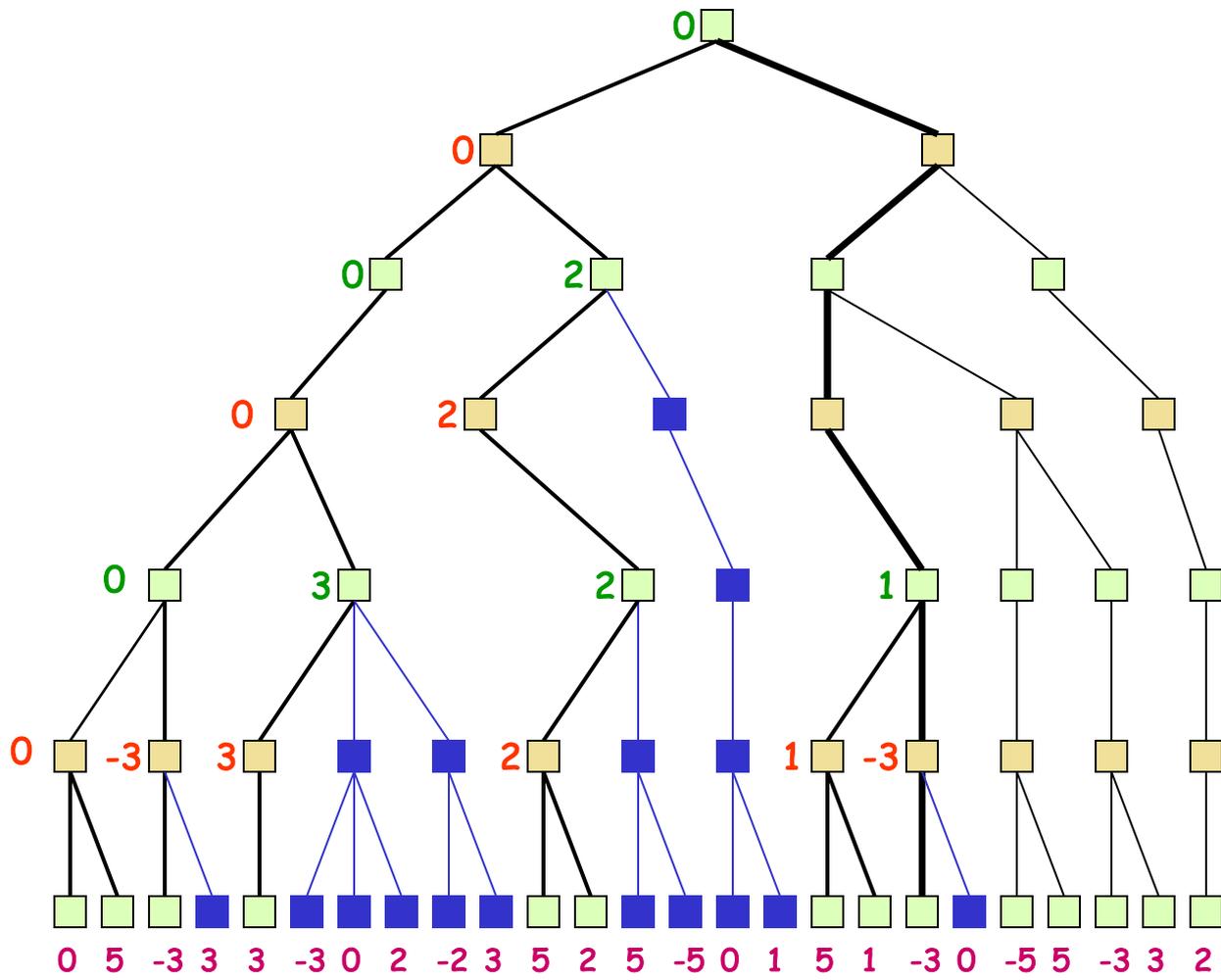


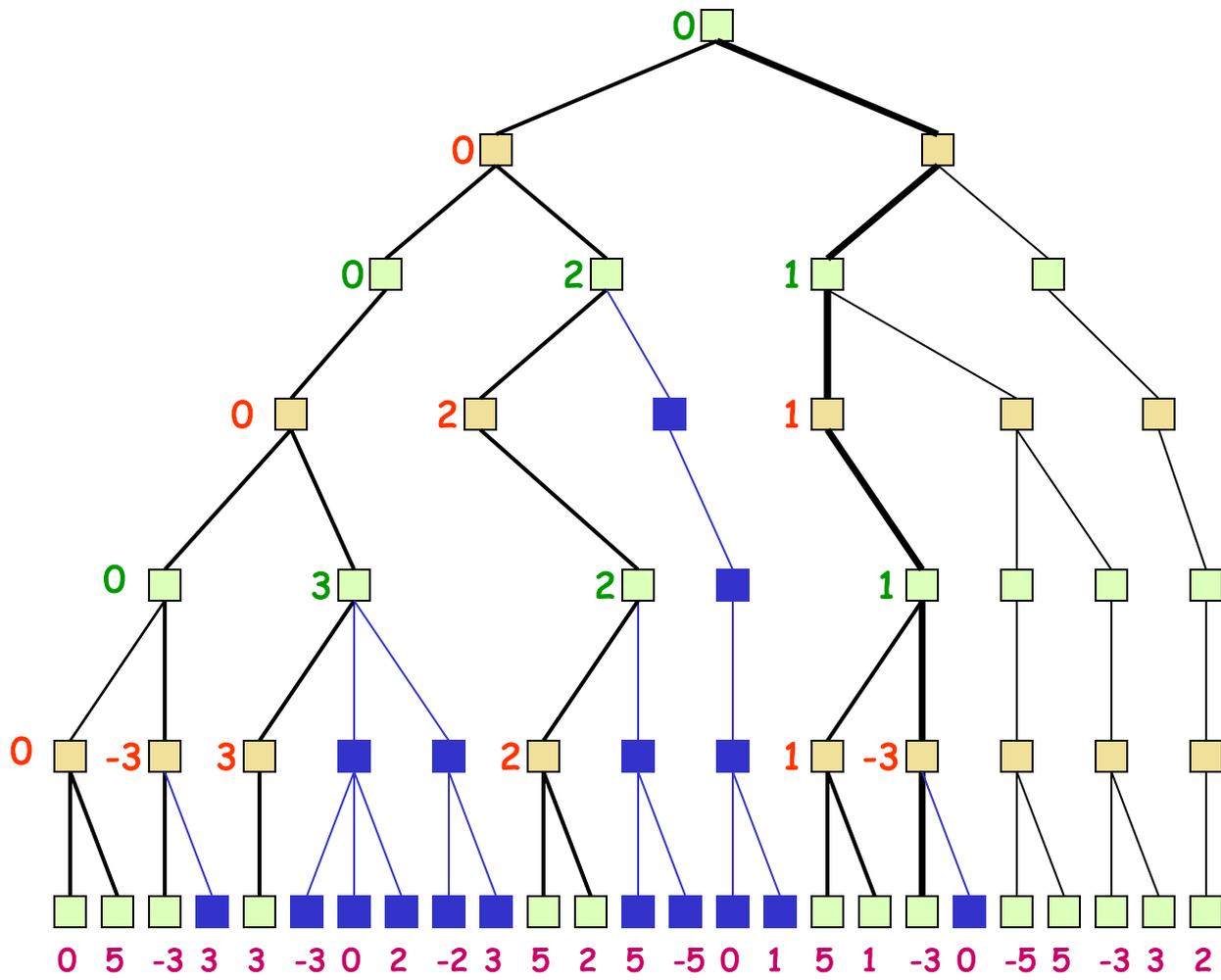


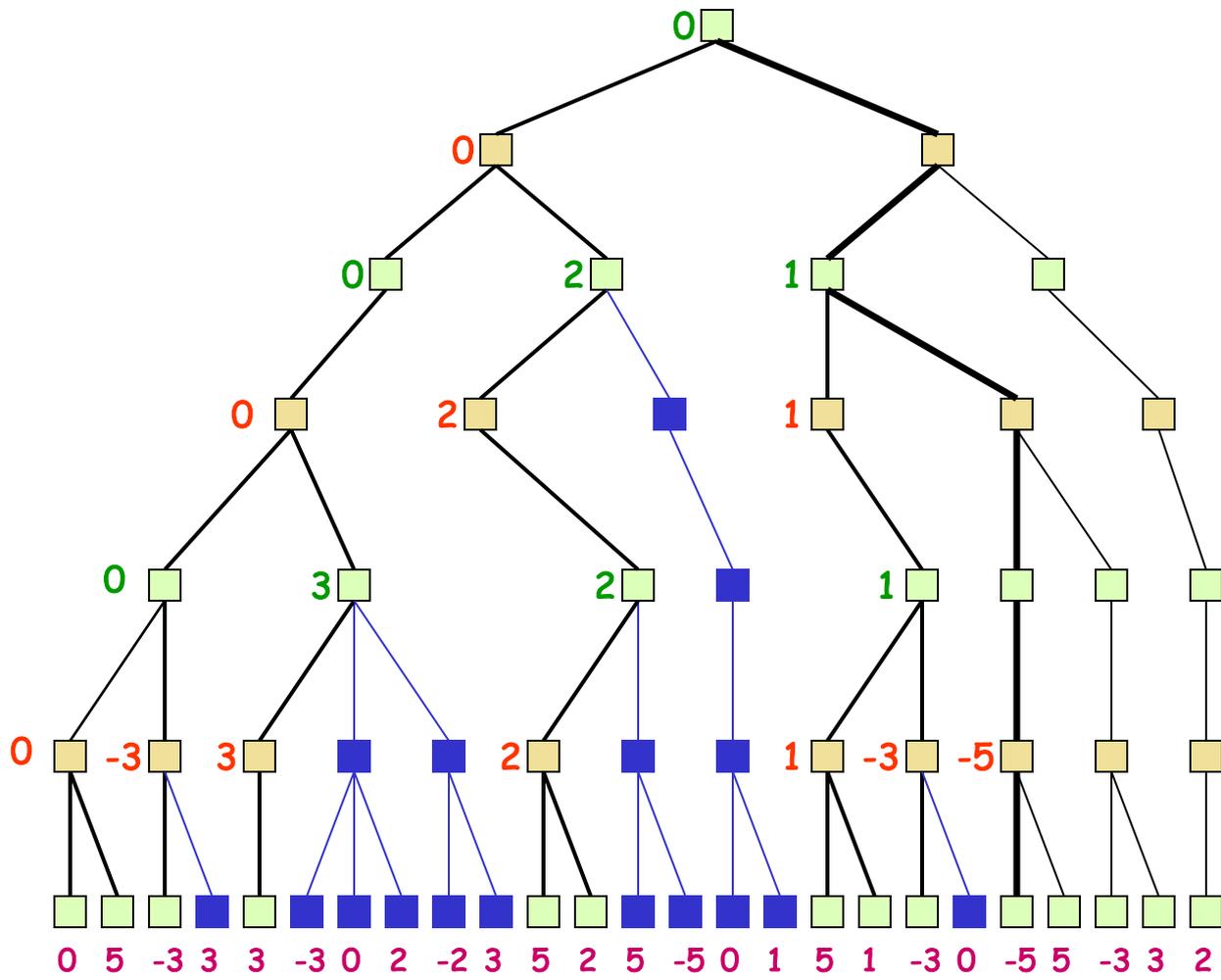


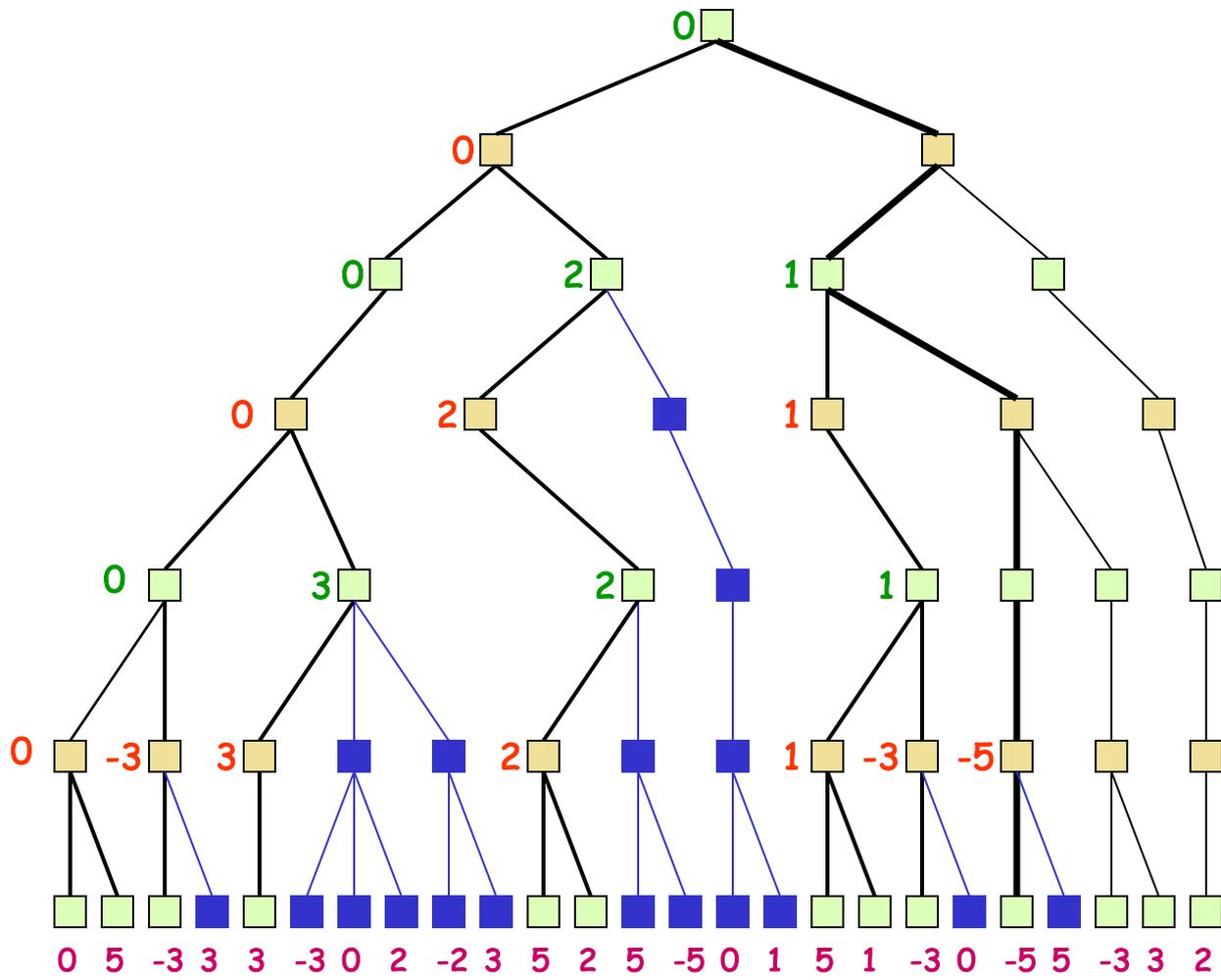


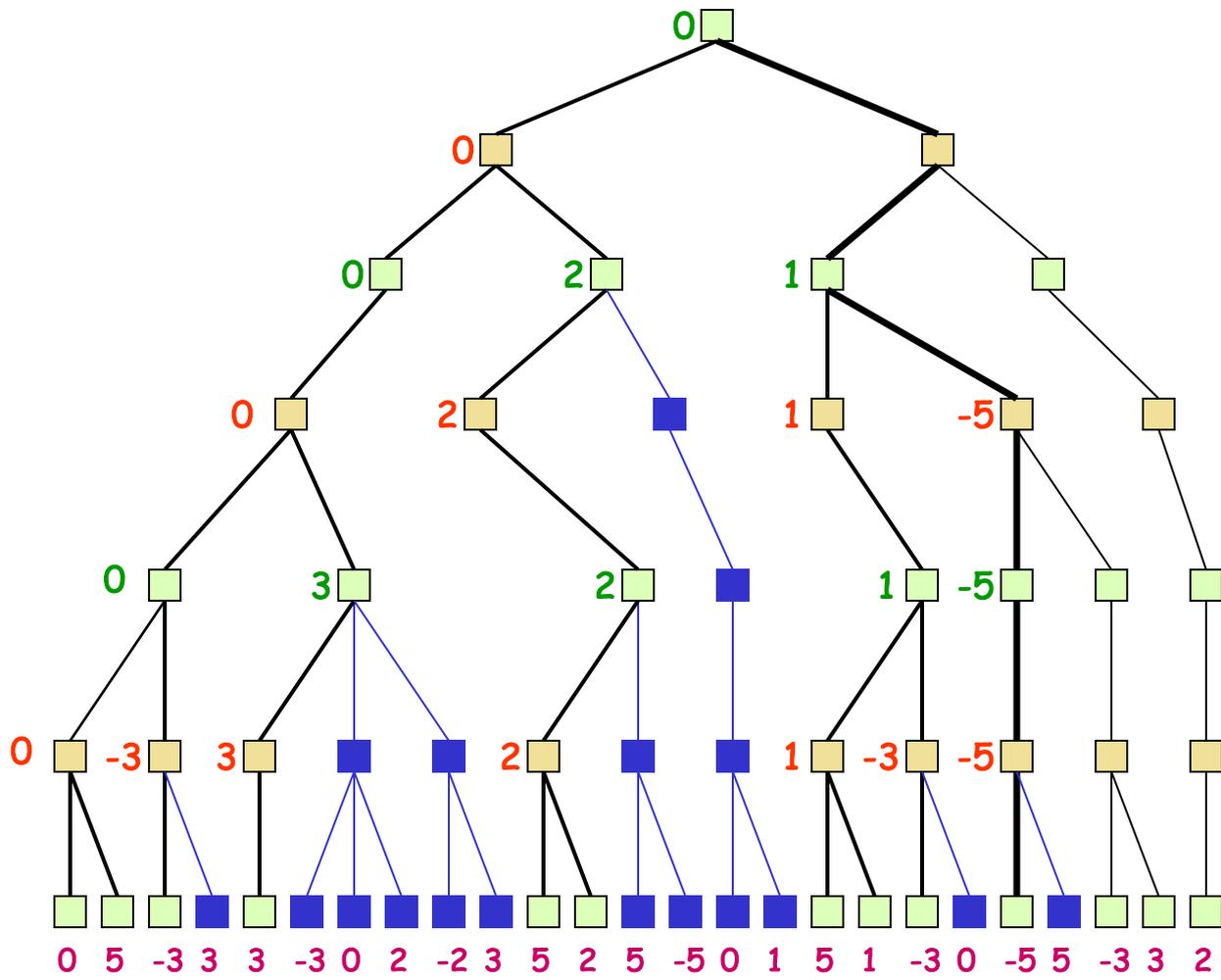


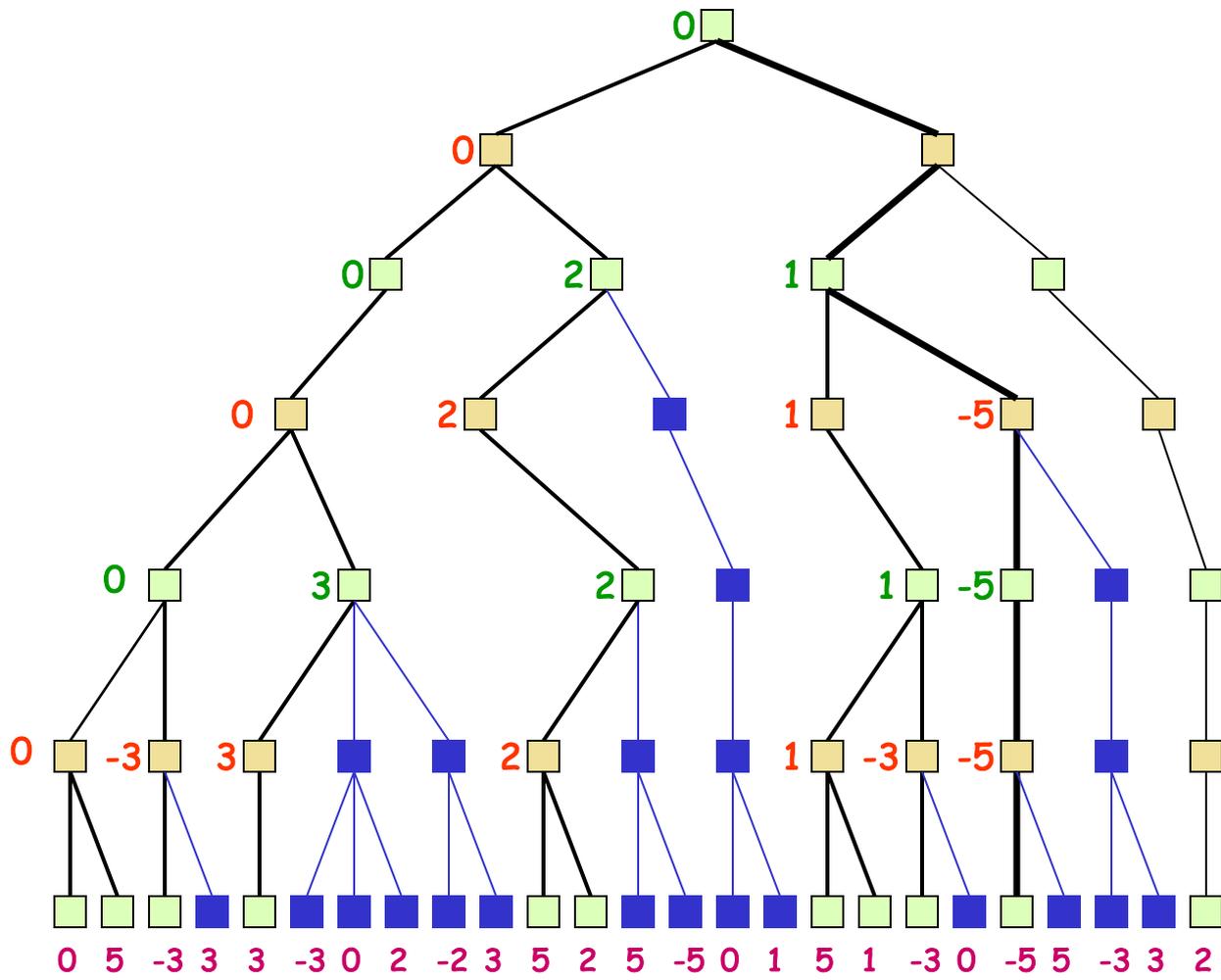


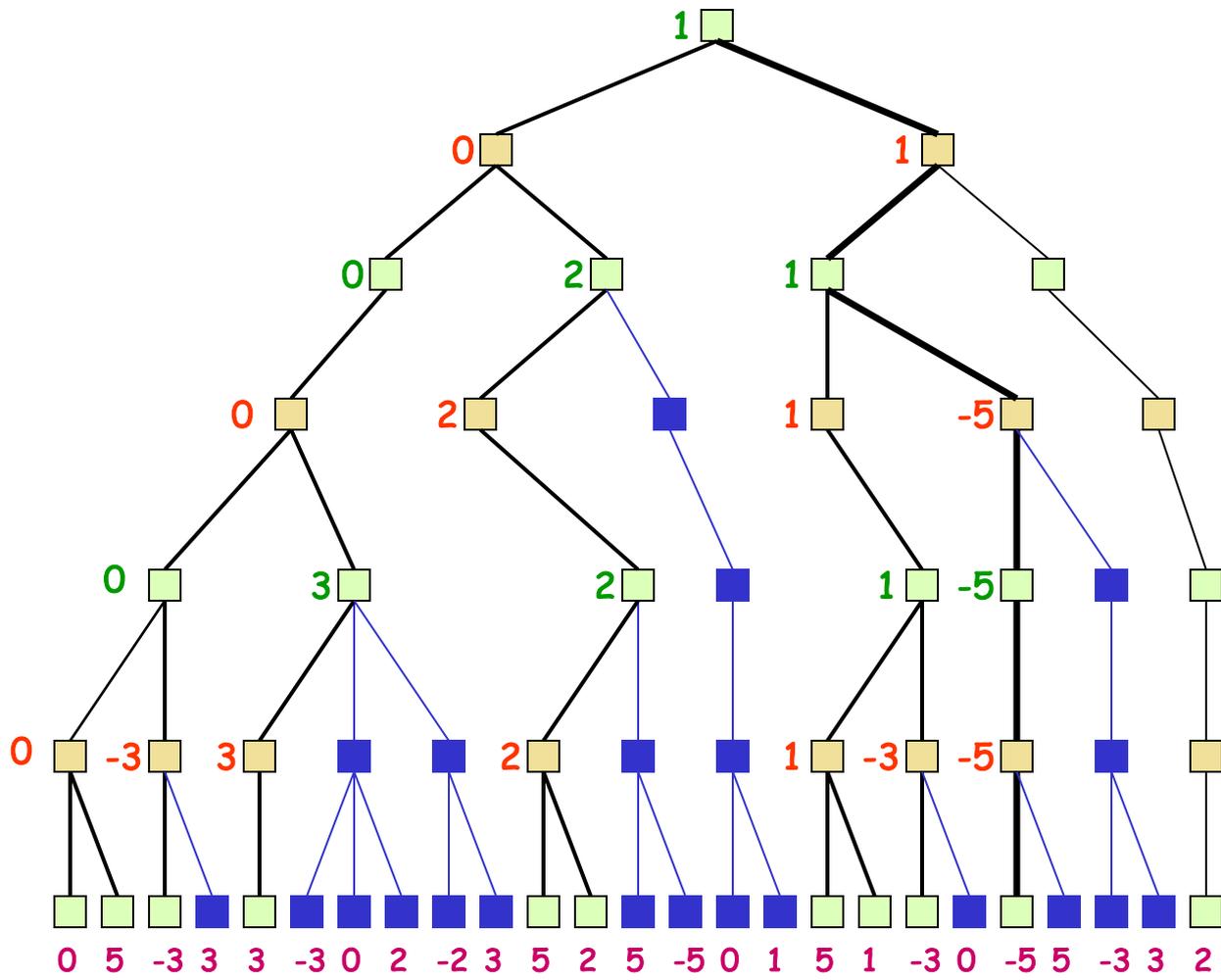


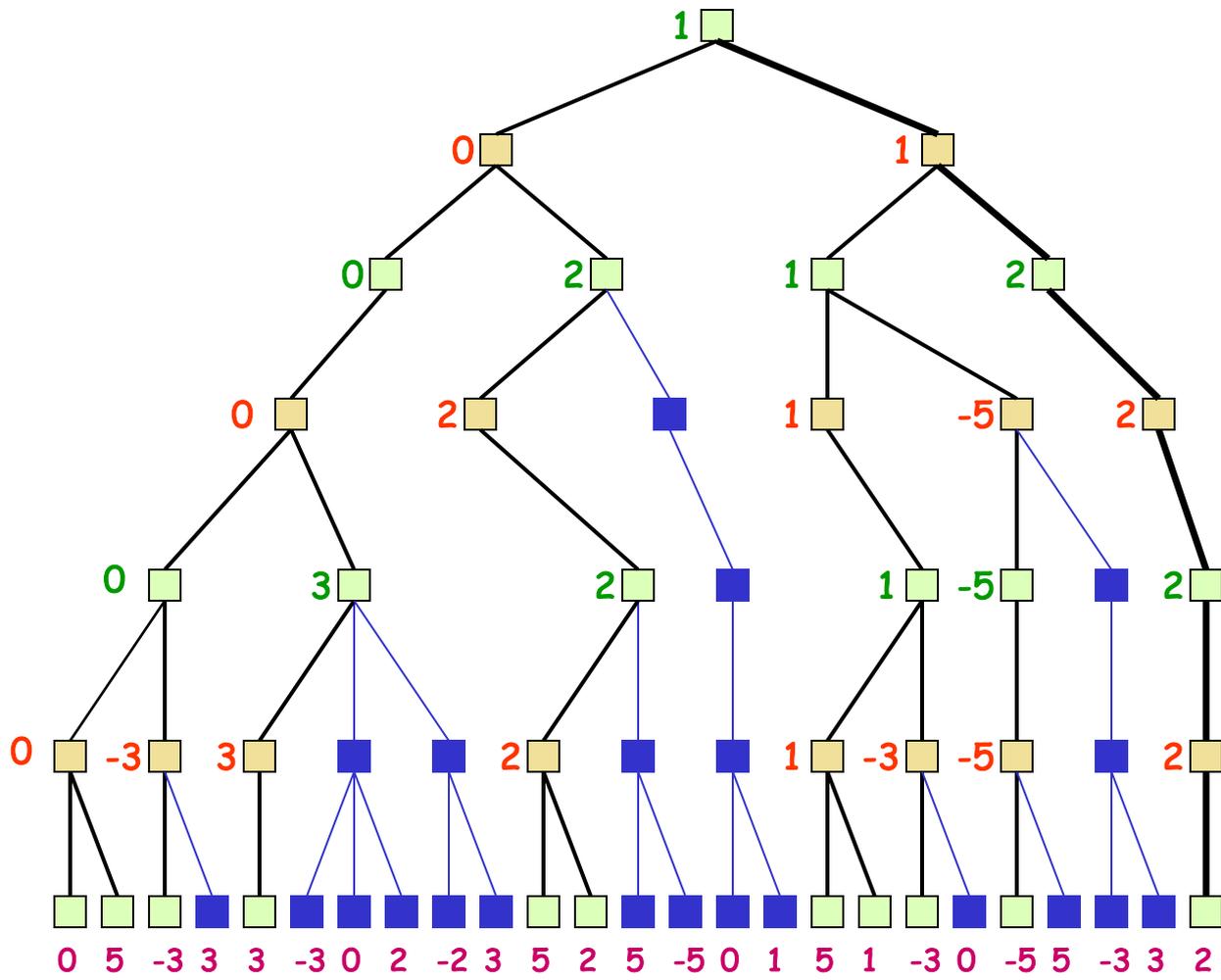




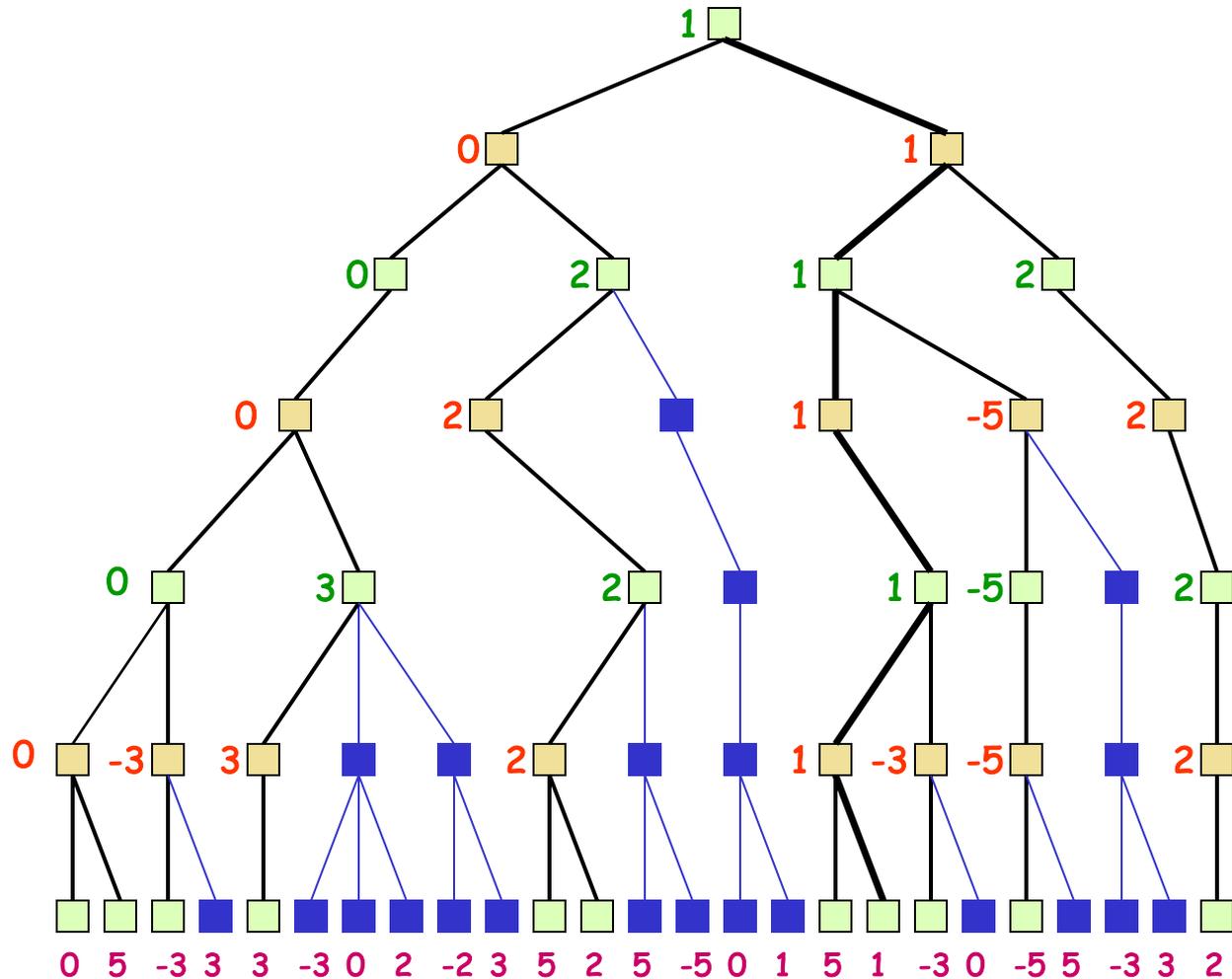








With alpha-beta we avoided computing a static evaluation metric for 14 of the 25 leaf nodes



Effectiveness of alpha-beta

- Alpha-beta guaranteed to compute same value for root node as minimax, but with \leq computation
- **Worst case:** no pruning, examine b^d leaf nodes, where nodes have b children & d -ply search is done
- **Best case:** examine only $(2b)^{d/2}$ leaf nodes
 - You can search twice as deep as minimax!
 - **Occurs if each player's best move is 1st alternative**
- In [Deep Blue](#), alpha-beta pruning reduced effective branching factor from ~ 35 to ~ 6

Many other improvements

- **Adaptive horizon + iterative deepening**
- **Extended search:** retain $k > 1$ best paths (not just one) extend tree at greater depth below their leaf nodes to help dealing with “horizon effect”
- **Singular extension:** If move is obviously better than others in node at horizon h , expand it
- Use [transposition tables](#) to deal with repeated states

Summary

- Simple 2-player, zero-sum, deterministic, perfect information games are popular and let us explore adversarial search
- Use a static evaluator and look ahead to choose move
- Computing static evaluator uses most computing
- Minimax makes best choice for next move
- Alpha-beta gives same answer, but often with much less work