

## Quadric Error Metrics

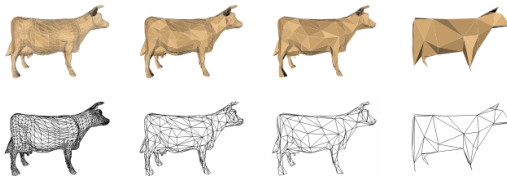


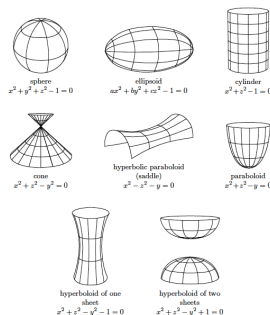
Figure 2.7: Fixed set of levels of detail for a cow model.

### To do

- Continue to work on ray programming assignment
- Continue to think about final project
- Next Monday: Dr. Lee Butler's guest lecture

### Background: quadrics

- A quadrics are all surfaces that can be expressed as a second degree polynomial in  $x$ ,  $y$ , and  $z$ .



### Quadric surfaces

- General implicit form

$$ax^2 + 2bxy + 2cxz + 2dxw + ey^2 + 2fyz + 2gyw + hz^2 + 2izw + jw^2 = 0.$$

- Matrix form

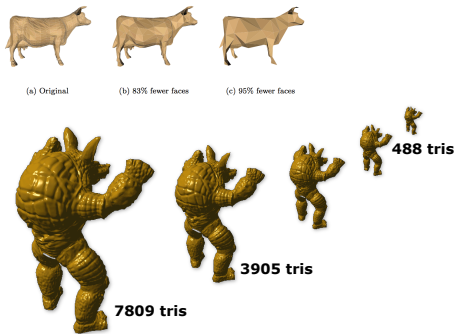
– Setting  $w$  to 1, this provides the ability to position the quadrics in space

$$\mathbf{x}^t Q \mathbf{x} = 0,$$

where  $\mathbf{x} = \begin{bmatrix} x \\ y \\ z \\ w \end{bmatrix}$ ,  $\mathbf{x}^t = [x \ y \ z \ w]$ , and  $Q = \begin{bmatrix} a & b & c & d \\ b & e & f & g \\ c & f & h & i \\ d & g & i & j \end{bmatrix}$

## Surface Simplification: Goals

- Efficiency (70000 to 100 faces in 15s in 1997)
- High quality, feature preserving (primary appearance emphasized rather than topology)
- Generally, non-manifold models, collapse disjoint regions



## Background: Manifold

- Mathematical term of a surface, all of those points have a neighborhood which is topologically equivalent to a disk.
- A manifold with boundary is a surface all of whose points have either a disk or a half-disk.
- A polygonal surface is a manifold (with boundary) if every edge has exactly two incident faces, and the neighborhood of every vertex consists of a closed loop of faces (or a single fan of faces on the boundary)

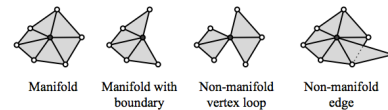
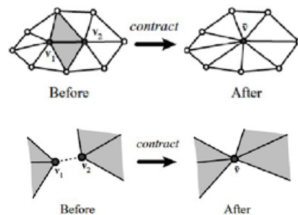


Figure 2.1: Neighborhoods of a given vertex (in black). On the left, two manifold neighborhoods. On the right, two non-manifold neighborhoods.

## Overview and Resources

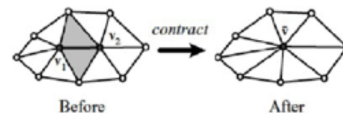
- Garland and Heckbert SIGGRAPH 97 paper
  - Greedy decimation algorithm
  - Pair collapse (allow edge + non-edge collapse)
- Evaluate potential collapse
- Determine optimal new vertex locations



- Garland website, implementation notes in his thesis  
<http://mgarland.org/research/quadrics.html>  
<http://mgarland.org/research/thesis.html>
- Notes in this and the previous lecture

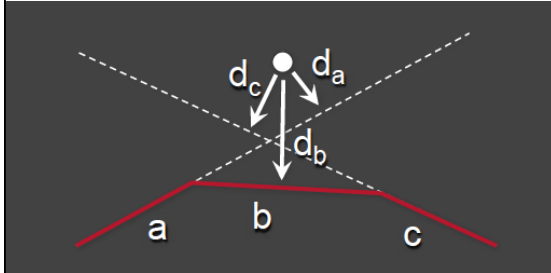
## Algorithm Outline

- Restrict process to a set of *valid pairs*:
  - $(v_i, v_j)$  is an edge, or
  - $\|v_i - v_j\| < t$ , a threshold
    - $t = 0$  restricts to edge contraction
    - $t \gg 0$  can connect distant regions or yield  $O(n^2)$  pairs
- Iteratively remove *best* pair and update valid pairs list:
  - Each vertex has a set with the pairs it belongs to:
    - $v_i \mapsto \text{Pairs}(v_i)$
    - $(v_i, v_j) \rightarrow \bar{v} \Rightarrow \text{Pairs}(\bar{v}) = \text{Pairs}(v_i) \cup \text{Pairs}(v_j)$
- But how to choose *best* pair?



### Choose best pair

- Based on point-to-plane distance
- Better quality than point-to-point



### Background: Computing Planes

- Each triangle in mesh has associated plane

$$ax + by + cz + d = 0$$

- For a triangle, find its (normalized) normal using cross products

$$\vec{n} = \frac{\vec{AB} \times \vec{AC}}{|\vec{AB} \times \vec{AC}|} \quad \vec{n} \cdot \vec{v} - \vec{A} \cdot \vec{v} = 0$$

- Plane equation:

$$\vec{n} = \begin{pmatrix} a \\ b \\ c \end{pmatrix} \quad d = -\vec{A} \cdot \vec{v}$$

### Quadric Error Metrics

- Sum of squared distances from vertex to planes

$$\Delta_v = \sum_p \text{Dist}(\mathbf{v}, \mathbf{p})^2 \quad \Delta = \sum_p (\mathbf{p}^T \mathbf{v})^2$$

$$\mathbf{v} = \begin{pmatrix} x \\ y \\ z \\ 1 \end{pmatrix}, \quad \mathbf{p} = \begin{pmatrix} a \\ b \\ c \\ d \end{pmatrix}$$

$$\text{Dist}(\mathbf{v}, \mathbf{p}) = ax + by + cz + d = \mathbf{p}^T \mathbf{v} \quad = \sum_p \mathbf{v}^T \mathbf{p} \mathbf{p}^T \mathbf{v}$$

$$= \mathbf{v}^T \left( \sum_p \mathbf{p} \mathbf{p}^T \right) \mathbf{v} = \mathbf{v}^T \mathbf{Q} \mathbf{v}$$

- Common mathematical trick:  
Quadratic form = symmetric matrix Q multiplied twice by a vector.

### Quadric Error Metrics

- A 4x4 symmetric matrix
- Storage efficient: 10 floating point number per vertex
- Initially, error is 0 for all vertices

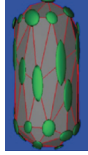
$$\mathbf{v}^T \mathbf{Q} \mathbf{v} = \begin{bmatrix} x & y & z & 1 \end{bmatrix} \begin{bmatrix} q_{11} & q_{12} & q_{13} & q_{14} \\ q_{12} & q_{22} & q_{23} & q_{24} \\ q_{13} & q_{23} & q_{33} & q_{34} \\ q_{14} & q_{24} & q_{34} & q_{44} \end{bmatrix} \begin{bmatrix} x \\ y \\ z \\ 1 \end{bmatrix}$$

- 2<sup>nd</sup> degree polynomial in x, y, and z
- Level surface ( $\mathbf{v}^T \mathbf{Q} \mathbf{v} = k$ ) is a quadric surface
  - Ellipsoid, parabolic surface, hyperboloid, plane etc.

$$\mathbf{v}^T \mathbf{Q} \mathbf{v} = q_{11}x^2 + 2q_{12}xy + 2q_{13}xz + 2q_{14}x + q_{22}y^2 + 2q_{23}yz + 2q_{24}y + q_{33}z^2 + 2q_{34}z + q_{44}$$

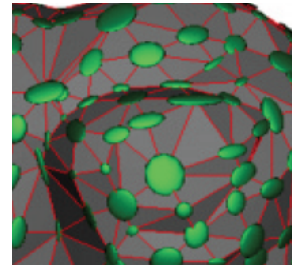
## But What Are These Quadrics Really Doing?

- Almost always ellipsoids
  - When  $Q$  is positive definite
- Characterize error at vertex
  - Vertex at center of each ellipsoid
  - Move it anywhere on ellipsoid with constant error
- Capture local shape of surface
  - Stretch in least curved direction



## Quadric Visualization

- Ellipsoids: iso-error surfaces
- Smaller ellipsoid = greater error for a given motion
- Lower error for motion parallel to surface
- Lower error in flat regions than at corners
- Elongated in “cylindrical” regions



## Using Quadrics

- Approximate error of edge collapses
  - Each vertex  $v$  has associated quadric  $Q$
  - Error of collapsing  $v_1$  and  $v_2$  to  $v'$  is  $v'^t Q_1 v' + v'^t Q_2 v'$
  - Quadric for new vertex  $v'$  is  $Q' = Q_1 + Q_2$

- Find optimal location  $v'$  after collapse

$$Q' = \begin{bmatrix} q_{11} & q_{12} & q_{13} & q_{14} \\ q_{12} & q_{22} & q_{23} & q_{24} \\ q_{13} & q_{23} & q_{33} & q_{34} \\ q_{14} & q_{24} & q_{34} & q_{44} \end{bmatrix}$$

$$\min_{v'} v'^t Q' v': \frac{\partial}{\partial x} = \frac{\partial}{\partial y} = \frac{\partial}{\partial z} = 0$$

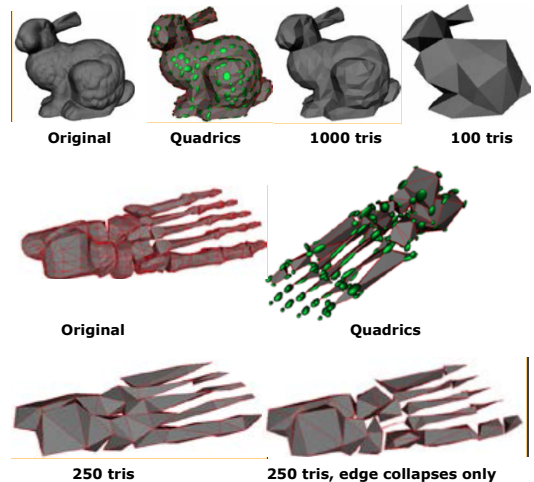
$$\begin{bmatrix} q_{11} & q_{12} & q_{13} & q_{14} \\ q_{12} & q_{22} & q_{23} & q_{24} \\ q_{13} & q_{23} & q_{33} & q_{34} \\ 0 & 0 & 0 & 1 \end{bmatrix} v' = \begin{bmatrix} 0 \\ 0 \\ 0 \\ 1 \end{bmatrix}$$

$$v' = \begin{bmatrix} q_{11} & q_{12} & q_{13} & q_{14} \\ q_{12} & q_{22} & q_{23} & q_{24} \\ q_{13} & q_{23} & q_{33} & q_{34} \\ 0 & 0 & 0 & 1 \end{bmatrix}^{-1} \begin{bmatrix} 0 \\ 0 \\ 0 \\ 1 \end{bmatrix}$$

## Algorithm Summary

- Compute the  $Q$  matrices for all the initial vertices
- Select all valid pairs
- Compute the optimal contraction target  $v$  for each valid pair. The error  $v^t(Q_1+Q_2)v'$  of this target vertex becomes the **cost** of contracting that pair.
- Place all pairs in a heap keyed on cost with minimum cost pair on the top
- Iteratively remove the least cost pair, contract this pair, and update the costs of all valid pairs of interest
  - Compute  $Q_i$  for all vertices  $v_i$
  - Determine valid pairs
  - Compute optimal contraction target and associated quadric error for each pair
  - Place all pairs in a heap, ordered by smallest error
  - Repeat
    - Get least error pair  $(v_i, v_j)$  from heap
    - Contract pair (move edges to  $v$ , remove degenerate planes)
    - Update cost for all pairs involving  $v_i$  and  $v_j$
  - Until done.

## Results



## Additional Details

- Preserving boundaries / discontinuities (weight quadrics by appropriate penalty factors)
- Preventing mesh inversion (flipping of orientation): compare normal of neighboring faces, before after
- Has been modified for many other applications
  - E.g., in silhouettes, want to make sure volume always increases, never decreases
  - Take color and texture into account (follow up paper)

## Implementation Tips

- Incremental, test, debug, simple cases
- Find good data structure for heap etc.
- May help to visualize error quadrics if possible
- Challenging, but hopefully rewarding assignment