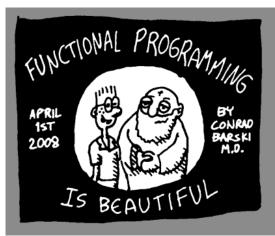
Functional Programming in Scheme and Lisp



http://www.lisperati.com/landoflisp/

Overview

- In a functional programming language, functions are first class objects
- You can create them, put them in data structures, compose them, specialize them, apply them to arguments, etc.
- We'll look at how functional programming things are done in Lisp

eval

- Remember: Lisp code is just an s-expression
- You can call Lisp's evaluation process with the eval function

```
> (define s (list 'cadr ' ' (one two three)))
> s
(cadr ' (one two three))
> (eval s)
two
> (eval (list 'cdr (car '((quote (a . b)) c))))
b
```

apply

• *apply* takes a function & a list of arguments for it & returns the result of applying the function to them

```
> (apply + '(1 2 3))
```

• It can be given any number of arguments, so long as the last is a list:

```
> (apply + 1 2 ' (3 4 5))
15
```

 A simple version of apply could be written as (define (apply f list) (eval (cons f list)))

lambda expression

• A *lambda expression* is a list of the symbol *lambda*, followed by a list of *parameters*, followed by a *body* of one or more expressions:

```
> (define f (lambda (x) (+ x 2)))
> f
###
> (f 100)
102
> ( (lambda (x) (+ x 2)) 100
102
```

lambda

- The define special form creates a function and gives it a name
- However, functions don't have to have names, and we don't need to use define to create them
- The primitive way to create functions is to use the *lambda* special form
- These are often called lambda expressions, e.g. (lambda (x) (+ x 1))

Lambda expression

- lambda is a special form
- When evaluated, it creates a function and returns a reference to it
- The function does not have a name
- A lambda expression can be the first element of a function call:
 - > ((lambda (x) (+ x 100)) 1) 101
- Other languages like python and javascript have adopted the idea

define vs. define

- The define special form comes in two varieties
- The three expressions to the right are entirely equivalent
- The first define form is just more familiar and convenient when defining a function

lambdas in other languages

• Lambda expressions are found in many modern languages, e.g., Python:

```
>>> f = lambda x,y: x*x + y
>>> f
<function <lambda> at 0x10048a230>
>>> f(2, 3)
7
>>> (lambda x,y: x*x+y)(2,3)
7
```

Functions as objects

 While many PLs allow functions as arguments, nameless lambda functions add flexibility

```
> (sort '((a 100)(b 10)(c 50))
(lambda (x y) (< (second x) (second y))))
((b 10) (c 50) (a 100))
```

There is no need to give the function a name

Mapping functions

- Lisp & Scheme have several mapping functions
- map (mapcar in Lisp) is the most useful
- It takes a function and ≥1 lists and returns a list of the results of applying the function to elements taken from each list

```
> (map abs '(3 -4 2 -5 -6))
(3 4 2 5 6)
> (map + '(1 2 3) (4 5 6))
(5 7 9)
> (map '(1 2 3) '(4 5 6) '(7 8 9))
(12 15 18)
```

More map examples

```
> (map cons '(a b c) '(1 2 3))
((a . 1) (b . 2) (c . 3))
> (map (lambda (x) (+ x 10)) '(1 2 3))
(11 12 13)
> (map + '(1 2 3) '(4 5))
map: all lists must have same size; arguments were:
#<procedure:+> (1 2 3) (4 5)
=== context ===
/Applications/PLT/collects/scheme/private/misc.ss:
74:7
```

Defining map

```
Defining a simple "one argument" version of map is easy

(define (map1 func list)

(if (null? list)

null

(cons (func (first list))

(map1 func (rest list)))))
```

Define Lisp's every and some

- every and some take a predicate and one or more sequences
- When given just one sequence, they test whether the elements satisfy the predicate

```
> (every odd? '(1 3 5))
#t
> (some even? '(1 2 3))
#t
```

• If given >1 sequences, the predicate takes as many args as there are sequences and args are drawn one at a time from them:

```
> (every > '(1 3 5) '(0 2 4)) #t
```

Defining every is easy

Define some similarly

filter

```
(filter <f>  returns a list of the elements of  which satisfy the predicate <f> (filter odd? '(0 1 2 3 4 5)) (1 3 5) (101.1 98.6 98.1 99.4 102.2)) (101.1 99.4 102.2)
```

Will this work?

- You can prove that P is true for some list element by showing that it isn't false for every one
- Will this work?

```
> (define (some1 f list)
     (not (every1 (lambda (x) (not (f x))) list)))
> (some1 odd? '(2 4 6 7 8))
#t
> (some1 (lambda (x) (> x 10)) '(4 8 10 12))
#t
```

Example: filter

Example: filter

• Define *integers* as a function that returns a list of integers between a min and max

```
(define (integers min max)
  (if (> min max)
    null
    (cons min (integers (add1 min) max))))
```

• Do prime? as a predicate that is true of prime numbers and false otherwise

```
> (filter prime? (integers 2 20) )
(2 3 5 7 11 13 17 19)
```

Here's another pattern

• We often want to do something like sum the elements of a sequence

```
(define (sum-list I)

(if (null? I)

0

(+ (first I) (sum-list (rest I)))))
```

• Other times we want their product

```
(define (multiply-list I)
    (if (null? I)
    1
    (* (first I) (multiply-list (rest I)))))
```

Here's another pattern

• We often want to do something like sum the elements of a sequence

```
(define (sum-list I)
    (if (null? I)
     0
     (+ (first I) (sum-list (rest I)))))
```

• Other times we want their product

```
(define (multiply-list I)
  (if (null? I)
    1
    (* (first I) (multiply-list (rest I)))))
```

Example: reduce

 Reduce takes (i) a function, (ii) a final value and (iii) a list of arguments

```
Reduce of +, 0, (v1 \ v2 \ v3 \dots vn) is just V1 + V2 + V3 + \dots Vn + 0
```

• In Scheme/Lisp notation:

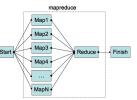
```
> (reduce + 0 '(1 2 3 4 5))
15
(reduce * 1 '(1 2 3 4 5))
120
```

Example: reduce

```
(define (reduce function final list)
  (if (null? list)
    final
      (function
          (first list)
          (reduce function final (rest list)))))
```

The roots of mapReduce

 MapReduce is a software framework developed by Google for parallel computation on large datasets on computer clusters



hedooo

- It's become an important way to exploit parallel computing using conventional programming languages and techniques.
- See Apache's <u>Hadoop</u> for an open source version
- The framework was inspired by functional programming's map, reduce & side-effect free programs

```
(define (sum-list list)

;; returns the sum of the list elements
(reduce + 0 list))

(define (mul-list list)

;; returns the sum of the list elements
(reduce * 1 list))

(define (copy-list list)

;; copies the top level of a list
(reduce cons '() list))

(define (append-list list)

;; appends all of the sublists in a list
(reduce append '() list))
```

Function composition

- Math notation: g •h is a composition of functions g and h
- If $f=q \bullet h$ then f(x)=g(h(x))
- Composing functions is easy in Scheme

```
> compose
###compose>
> (define (sq x) (* x x))
> (define (dub x) (* x 2))
> (sq (dub 10))
400
> (dub (sq 10))
200
```

```
> (define sd (compose sq
dub))
> (sd 10)
400
> ((compose dub sq) 10)
200
```

Defining compose

- Here's compose for two functions in Scheme (define (compose2 f g) (lambda (x) (f (g x))))
- Note that compose calls lambda which returns a new function that applies f to the result of applying g to x
- We'll look at how the variable environments work to support this in the next topic, closures
- But first, let's see how to define a general version of compose taking any number of args

Functions with any number of args

 Defining functions that takes any number of arguments is easy in Scheme

```
(define (foo . args) (printf "My args: ~a\n" args)))
```

 If the parameter list ends in a symbol as opposed to null (cf. dotted pair), then it's value is the list of the remaining arguments' values

```
(define (f x y . more-args) ...)
(define (map f . lists) ...)
```

Compose in Scheme

A general every

 We can easily re-define other functions to take more than one argument

- (every > '(1 2 3) '(0 2 3)) => #t
- (every > '(1 2 3) '(0 20 3)) => #f

Functional Programming Summary

- List is the archetypal functional programming language
- It treated functions as first-class objects and uses the same representation for data & code
- The FP paradigm is a good match for many problems, esp. ones involving reasoning about or optimizing code or parallel execution
- While no pure FP languages are considered mainstream, many PLs support a FP style