# Functional Programming in Scheme and Lisp



#### Overview

- In a functional programming language, functions are first class objects
- You can create them, put them in data structures, compose them, specialize them, apply them to arguments, etc.
- We'll look at how functional programming things are done in Lisp

#### eval

- Remember: Lisp code is just an s-expression
- You can call Lisp's evaluation process with the eval function
  - > (define s (list 'cadr '' (one two three)))
  - > s
  - (cadr ' (one two three))
  - > (eval s)
  - two
  - > (eval (list 'cdr (car '((quote (a . b)) c))))

b

# apply

 apply takes a function & a list of arguments for it & returns the result of applying the function to them > (apply + '(1 2 3))

• It can be given any number of arguments, so long as the last is a list:

>(apply + 1 2 '(3 4 5)) 15

 A simple version of apply could be written as (define (apply f list) (eval (cons f list)))

# lambda

- The *define* special form creates a function and gives it a name
- However, functions don't have to have names, and we don't need to use *define* to create them
- The primitive way to create functions is to use the *lambda* special form
- These are often called lambda expressions, e.g. (lambda (x) (+ x 1))

#### lambda expression

A lambda expression is a list of the symbol lambda, followed by a list of parameters, followed by a body of one or more expressions:
> (define f (lambda (x) (+ x 2)))
> f
#<proceedure:f>
> (f 100)
102
> ( (lambda (x) (+ x 2)) 100)
102

#### Lambda expression

- · lambda is a special form
- When evaluated, it creates a function and returns a reference to it
- The function does not have a name
- A lambda expression can be the first element of a function call:

>((lambda (x) (+ x 100)) 1) 101

• Other languages like python and javascript have adopted the idea

# define vs. define (define (add2 x) (+ x 2)) • The define spectrum of two x

(define add2

(lambda (x) (+ x 2)))

(define add2 #f) (set! add2 (lambda (x) (+ x 2)))

- The define special form comes in two varieties
- The three expressions to the right are entirely equivalent
- The first define form is just more familiar and convenient when defining a function

# Functions as objects

- While many PLs allow functions as arguments, nameless lambda functions add flexibility

  - ((b 10) (c 50) (a 100))
- There is no need to give the function a name

# lambdas in other languages

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Lambda expressions are found in many
modern languages, e.g., Python:
>> f = lambda x,y: x*x + y
>> f
<function <lambda> at 0x10048a230>
>>> f(2, 3)
7
>>> (lambda x,y: x*x+y)(2,3)
7
```

#### Mapping functions • Lisp & Scheme have several mapping functions • map (mapcar in Lisp) is the most useful • It takes a function and ≥1 lists and returns a list of the results of applying the function to elements taken from each list > (map abs '(3 -4 2 -5 -6)) (3 4 2 5 6) > (map + '(1 2 3) (4 5 6)) (5 7 9)

> (map + '(1 2 3) '(4 5 6) '(7 8 9))

(12 15 18)

#### More map examples

> (map cons '(a b c) '(1 2 3))
((a . 1) (b . 2) (c . 3))
> (map (lambda (x) (+ x 10)) '(1 2 3))
(11 12 13)
> (map + '(1 2 3) '(4 5))
map: all lists must have same size; arguments were:
#<procedure:+> (1 2 3) (4 5)

/Applications/PLT/collects/scheme/private/misc.ss:7

#### **Defining map**

Defining a simple "one argument" version of map is easy (define (map1 func list)

(if (null? list) null

(cons (func (first list))

(map1 func (rest list)))))

#### Define Lisp's every and some

- every and some take a predicate and one or more sequences
- When given just one sequence, they test whether the elements satisfy the predicate > (every odd? '(1 3 5)) #t
  - >(some even? '(1 2 3))

=== context ===

- #t
- If given >1 sequences, the predicate takes as many args as there are sequences and args are drawn one at a time from them:
   > (every > '(1 3 5) '(0 2 4)) #t

# Defining every is easy

(define (every1 f list)
;; note the use of the and function
(if (null? list)
 #t
 (and (f (first list))
 (every1 f (rest list))))))

#### **Define some similarly**

(define (some1 f list)
 (if (null? list)
 #f
 (or (f (first list))
 (some1 f (rest list)))))

# Will this work?

- You can prove that P is true for some list element by showing that it isn't false for every one
- Will this work?
  - > (define (some1 f list)

(not (every1 (lambda (x) (not (f x))) list)))

> (some1 odd? '(2 4 6 7 8))

#t

> (some1 (lambda (x) (> x 10)) '(4 8 10 12))

#t

#### filter

(filter <f> <list>) returns a list of the elements of <list> which satisfy the predicate <f> > (filter odd? '(0 1 2 3 4 5))

(135)

> (filter (lambda (x) (> x 98.6)) (101.1 98.6 98.1 99.4 102.2))

(101.1 99.4 102.2)

# **Example: filter**

> (filter1 even? '(1 2 3 4 5 6 7)) (2 4 6)

#### **Example: filter**

• Define *integers* as a function that returns a list of integers between a min and max

(define (integers min max) (if (> min max)

null

- (cons min (integers (add1 min) max))))
- Do prime? as a predicate that is true of prime numbers and false otherwise

> (filter prime? (integers 2 20) )
(2 3 5 7 11 13 17 19)

#### Here's another pattern

 We often want to do something like sum the elements of a sequence (define (sum-list I)

(if (null? l)

0

(+ (first I) (sum-list (rest I)))))

Other times we want their product
 (define (multiply-list I)

(if (null? l)

(\* (first I) (multiply-list (rest I)))))



# Example: reduce

- Reduce takes (i) a function, (ii) a final value and (iii) a list of arguments Reduce of +, 0, (v1 v2 v3 ... vn) is just V1 + V2 + V3 + ... Vn + 0
  In Scheme/Lisp notation:
  - > (reduce + 0 '(1 2 3 4 5))
  - 15

(reduce \* 1 '(1 2 3 4 5))

120

# **Example: reduce**

(define (reduce function final list) (if (null? list) final (function (first list) (reduce function final (rest list)))))

	Using reduce
(define (sum-list list)	Using reduce
;; returns the sum of the list elements	
(reduce + 0 list))	
(define (mul-list list)	
;; returns the sum of the list elen	nents
(reduce * 1 list))	
(define (copy-list list)	
;; copies the top level of a list	
(reduce cons '() list))	
(define (append-list list)	
;; appends all of the sublists in a	list
(reduce append '() list))	

# The roots of mapReduce

 <u>MapReduce</u> is a software framework developed by Google for parallel computation on large datasets on computer clusters



- It's become an important way to exploit parallel computing using conventional programming languages and techniques.
- See Apache's <u>Hadoop</u> for an open source version



 The framework was inspired by functional programming's map, reduce & side-effect free programs



- Math notation: *g* •*h* is a <u>composition</u> of functions *g* and *h*
- If *f*=*g* •*h* then *f*(*x*)=*g*(*h*(*x*)
- · Composing functions is easy in Scheme

> compose

- #<procedure:compose>
  > (define (sq x) (\* x x))
- > (define (dub x) (\* x 2))
- > (sq (dub 10))
- 400

> (dub (sq 10)) 200 > (define sd (compose sq dub)) > (sd 10) 400 > ((compose dub sq) 10) 200

# **Defining compose**

- Here's compose for two functions in Scheme (define (compose2 f g) (lambda (x) (f (g x))))
- Note that compose calls lambda which returns a new function that applies *f* to the result of applying *g* to *x*
- We'll look at how the variable environments work to support this in the next topic, closures
- But first, let's see how to define a general version of compose taking any number of args

# Functions with any number of args

• Defining functions that takes any number of arguments is easy in Scheme

(define (foo . args) (<u>printf</u> "My args: ~a\n" args)))

 If the parameter list ends in a symbol as opposed to null (cf. dotted pair), then it's value is the list of the remaining arguments' values

(define (f x y . more-args) ...)

(define (map f. lists) ... )

# **Compose in Scheme**

# A general every

- We can easily re-define other functions to take more than one argument

   (define (every fn . args)
   (cond ((null? args) #f)
   ((null? (first args)) #t)
   ((apply fn (map first args)))
   (apply every fn (map rest args)))
   (else #f)))
   (every > '(1 2 3) '(0 2 3)) => #t
- (every > '(1 2 3) '(0 20 3)) => #f

# **Functional Programming Summary**

- Lisp is the archetypal functional programming language
- It treated functions as first-class objects and uses the same representation for data & code
- The FP paradigm is a good match for many problems, esp. ones involving reasoning about or optimizing code or parallel execution
- While no pure FP languages are considered mainstream, many PLs support a FP style