

# Scheme in Python



## Scheme in Python

- We'll follow the approach taken in the scheme in scheme interpreter for scheme in Python
  - <http://cs.umbc.edu/courses/331/fall10/code/scheme/mcscheme/>
- While is similar to that used by Peter Norvig
  - <http://cs.umbc.edu/courses/331/fall10/code/python/scheme/>
- It's a subset of scheme with some additional limitations
- We'll also look at some extensions

## Key parts

- **S-expression** representation
- **parsing** input into s-expressions
- **Print** s-expression, serializing as text that can be read by (read)
- **Environment** representation plus defining, setting and looking up variables
- **Function** representation
- **Evaluation** of an s-expression
- **REPL** (read, eval, print) loop

## S-expression

- Scheme numbers are Python numbers
- Scheme symbols are Python strings
- Scheme strings are not supported, so simplify the parsing
- Scheme lists are python lists
  - No dotted pairs
  - No structure sharing or circular list
- #t and #f are True and False
- Other Scheme data types aren't supported

## Parse: string=>s-expression

```
def read(s):
    "Read a Scheme expression from a string."
    return read_from(tokenize(s))

parse = read

def tokenize(s):
    "Convert a string into a list of tokens."
    return s.replace("'", ' ').replace(')', ')').\
        replace('\n', ' ').strip().split()
```

## Parse: string=>s-expression

```
def read_from(tokens):
    "Read an expression from a sequence of tokens."
    if len(tokens) == 0:
        raise SchemeError('EOF while reading')
    token = tokens.pop(0)
    if '(' == token:
        L = []
        while tokens[0] != ')':
            L.append(read_from(tokens))
        tokens.pop(0)                      # pop off ')'
        return L
    elif ')' == token:
        raise SchemeError('unexpected )')
    else:
        return atom(token)
```

## Making atoms, loading files

```
def atom(token):
    "Numbers become numbers; every other token is a symbol."
    try: return int(token)
    except ValueError:
        try: return float(token)
        except ValueError:
            return Symbol(token)

def load(filename):
    "Reads expressions from file and evaluates them, returns void."
    tokens = tokenize(open(filename).read())
    while tokens:
        eval(read_from(tokens))
```

## print reverses read

```
def to_string(exp):
    "Convert a Python object back into a Lisp-readable string."
    if isa(exp, list):
        return '(' + ''.join(map(to_string, exp)) + ')'
    else:
        return str(exp)
```

## Environments

- An environment will just be a Python dictionary of symbol:value pairs
- Plus an extra key called **outer** whose value is the enclosing (parent) environment

## Env class

```
class Env(dict):
    "An environment: a dict of {var:val} pairs, with an outer Env."
    def __init__(self, parms=(), args=(), outer=None):
        self.update(zip(parms,args))
        self.outer = outer
    def find(self, var):
        "Returns the innermost Env where var appears."
        if var in self: return self
        elif self.outer: return self.outer.find(var)
        else: raise SchemeError("unbound variable " + var)
    def set(self, var, val): self.find(var)[var] = val
    def define(self, var, val): self[var] = val
    def lookup(self, var): return self.find(var)[var]
```

## Eval

```
def eval(x, env=global_env):
    "Evaluate an expression in an environment."
    if isa(x, Symbol): return env.lookup(x)
    elif not isa(x, list): return x
    elif x[0] == 'quote': return x[1]
    elif x[0] == 'if': return eval((x[2] if eval(x[1], env) else x[3]), env)
    elif x[0] == 'set!': env.set(x[1], eval(x[2], env))
    elif x[0] == 'define': env.define(x[1], eval(x[2], env))
    elif x[0] == 'lambda': return lambda *args: eval(x[2], Env(x[1], args, env))
    elif x[0] == 'begin': return [eval(exp, env) for exp in x[1:]] [-1]
    else:
        exps = [eval(exp, env) for exp in x]
        proc = exps.pop(0)
        return proc(*exps)
```

## Representing functions

- This interpreter represents scheme functions as Python functions. Sort of.
- Consider evaluating
  - (define sum (lambda (x y) (+ x y)))
- This binds sum to the evaluation of
  - ['lambda', ['x', 'y'], ['+', 'x', 'y']]
- Which evaluates the Python expression
  - lambda \*args: eval(x[3], Env(x[2], args, env))
  - = <function <lambda> at 0x10048aed8>
- Which remembers values of x and env

## Calling a function

- Calling a built-in function
  - $(+ 1 2) \Rightarrow [+ 1 2]$
  - $\Rightarrow [\text{built-in function add}, 1, 2]$
  - Evaluates  $\text{add}(1, 2)$
- Calling a user defined function
  - $(\text{sum} 1 2) \Rightarrow [\text{sum}, 1, 2]$
  - $\Rightarrow [\text{function } \lambda \text{ at...}, 1, 2]$
  - $\Rightarrow [\text{function } \lambda \text{ at...}](1, 2)$
  - $\Rightarrow \text{eval}([+, x, y], \text{Env}(x, y), [1, 2], \text{env})$

## repl()

```
def repl(prompt='pscm> '):
    """A prompt-read-eval-print loop"""
    print "pscheme, type control-D to exit"
    while True:
        try:
            val = eval(parse(raw_input(prompt)))
            if val is not None: print to_string(val)
        except EOFError:
            print "Leaving pscm"
            break
        except SchemeError as e:
            print "SCM ERROR: ", e.args[0]
        except:
            print "ERROR: ", sys.exc_info()[0]
    # if called as a script, execute repl()
    if __name__ == "__main__": repl()
```

## Extensions

- Pscm.py has lots of shortcomings that can be addressed
- More data types (e.g., strings)
- A better scanner and parser
  - Macros
- Functions with variable number of args
- Tail recursion optimization

## Strings should be simple

- But adding strings breaks our simple approach to tokenization
- We added spaces around parentheses and then split on white space
- But strings can have spaces in them
  - ☹
- The solution is to use regular expressions to match any of the tokens
- While we are at it, we can add more token types, ; comments, etc.

## quasiquote

- Lisp and Scheme use a single quote char to make the following s-expression a literal
- The back-quote char (`) is like ' except that any elements in the following expression preceded by a , or ,@ are evaluated
- ,@ means the following expression should be a list that is “spliced” into its place

```
> 'foo          > (define y '(1 2 3))
  foo           > `(foo ,@y bar)
  > (define x 100)   (foo 1 2 3 bar)
  > `(foo ,x bar)
  (foo 100 bar)
```

## ; comments

- Lisp and Scheme treat all text between a semi-colon and the following newline as a comment
- The text is discarded as it is being read

## Big hairy regular expression

```
r'" \s* (@|[(")]|"(?:[\\"].[^\\"])*";.*|[^\s(";,;)]*)(.*)"'
```

- Whitespace
- The next token alternatives:
  - ,@ quasiquote ,@ token
  - [(")] single character tokens
  - "(?:[\\"].[^\\"])\*" string
  - ;.\* comment
  - [^\s(";,;)]\* atom
- The characters after the next token

## Reading from ports

```
class InPort(object):
    "An input port. Retains a line of chars."
    tokenizer = r"\s*(@|[(")]|"(?:[\\"].[^\\"])*";.*|[^\s(";,;)]*)(.*)"'
    def __init__(self, file):
        self.file = file; self.line = ""
    def next_token(self):
        "Return the next token, reading new text into line buffer if needed."
        while True:
            if self.line == "": self.line = self.file.readline()
            if self.line == "": return eof_object
            token, self.line = re.match(InPort.tokenizer, self.line).groups()
            if token != " and not token.startswith(';'):
                return token
```

## Tail recursion optimization

- We can add some tail recursion optimization by altering our eval() function
- Consider evaluating
  - $(\text{if } (> v 0) (\text{begin } 1 (\text{begin } 2 (\text{twice } (- v 1)))) 0)$
- In an environment where
  - $v=1$
  - $\text{twice} = (\lambda(x) (* 2 x))$

## Tail recursion optimization

- Here's a trace of the recursive calls to eval

```
⇒ eval(x=(if (> v 0) (begin 1 (begin 2 (twice (- v 1)))) 0), env={'v':1})
  ⇒ eval(x=(begin 1 (begin 2 (twice (- v 1)))), env={'v':1})
    ⇒ eval(x=(begin 2 (twice (- v 1)))), env={'v':1})
      ⇒ eval(x=(twice (- v 1))), env={'v':1})
        ⇒ eval(x=(* 2 y), env={'v':1})
          ⇐ 0
          ⇐ 0
          ⇐ 0
          ⇐ 0
          ⇐ 0
          ⇐ 0
```

- Eliminate recursive eval calls by setting x and env to required new values & iterate

## Tail recursion optimization

- Wrap the eval code in a loop, use return to exit, otherwise set x and env to new values
  - Here's a trace of the recursive calls to eval
- ```
⇒ eval(x=(if (> v 0) (begin 1 (begin 2 (twice (- v 1)))), env={'v':1})
  x = (begin 1 (begin 2 (twice (- v 1))))
  x = (begin 2 (twice (- v 1)))
  x = (twice (- v 1)))
  x = (* 2 y); env = {'y':0}
  ⇐ 0
```
- No recursion: faster and does not grow stack

## User defined functions

- We'll have to unpack our representation for user defined functions

- Define a simple class for a procedure

```
class Procedure(object):
    "A user-defined Scheme procedure."
    def __init__(self, parms, exp, env):
        self.parms, self.exp, self.env = parms, exp, env
    def __call__(self, *args):
        return eval(self.exp, Env(self.parms, args, self.env))
```

- Evaluating a lambda will just instantiate this

```

def eval(x, env=global_env):
    while True:
        if isa(x, Symbol): return env.lookup(x)
        elif not isa(x, list): return x
        elif x[0] == 'quote': return x[1]
        elif x[0] == 'if': x = x[2] if eval(x[1], env) else x[3]
        elif x[0] == 'set!':
            env.set(x[1], x[2])
            return None
        elif x[0] == 'define':
            env.define(x[1], x[2])
            return None
        elif x[0] == 'lambda': return Procedure(x[1], x[2], env)
        elif x[0] == 'begin':
            for exp in x[1:-1]: eval(exp, env)
            x = exp[-1]
        else:
            exps = [eval(exp, env) for exp in x]
            proc = exps.pop(0)
            if isa(proc, Procedure):
                x, env = proc.exp, Env(proc.params, exps, proc.env)
            else: return proc(*exps)

```

## Eval